

Logistics

- Grades for the midterm and the course average have been posted on learn.unm.edu.
- Reviews are due today at 6:00pm
- Project 3 will be posted this afternoon.

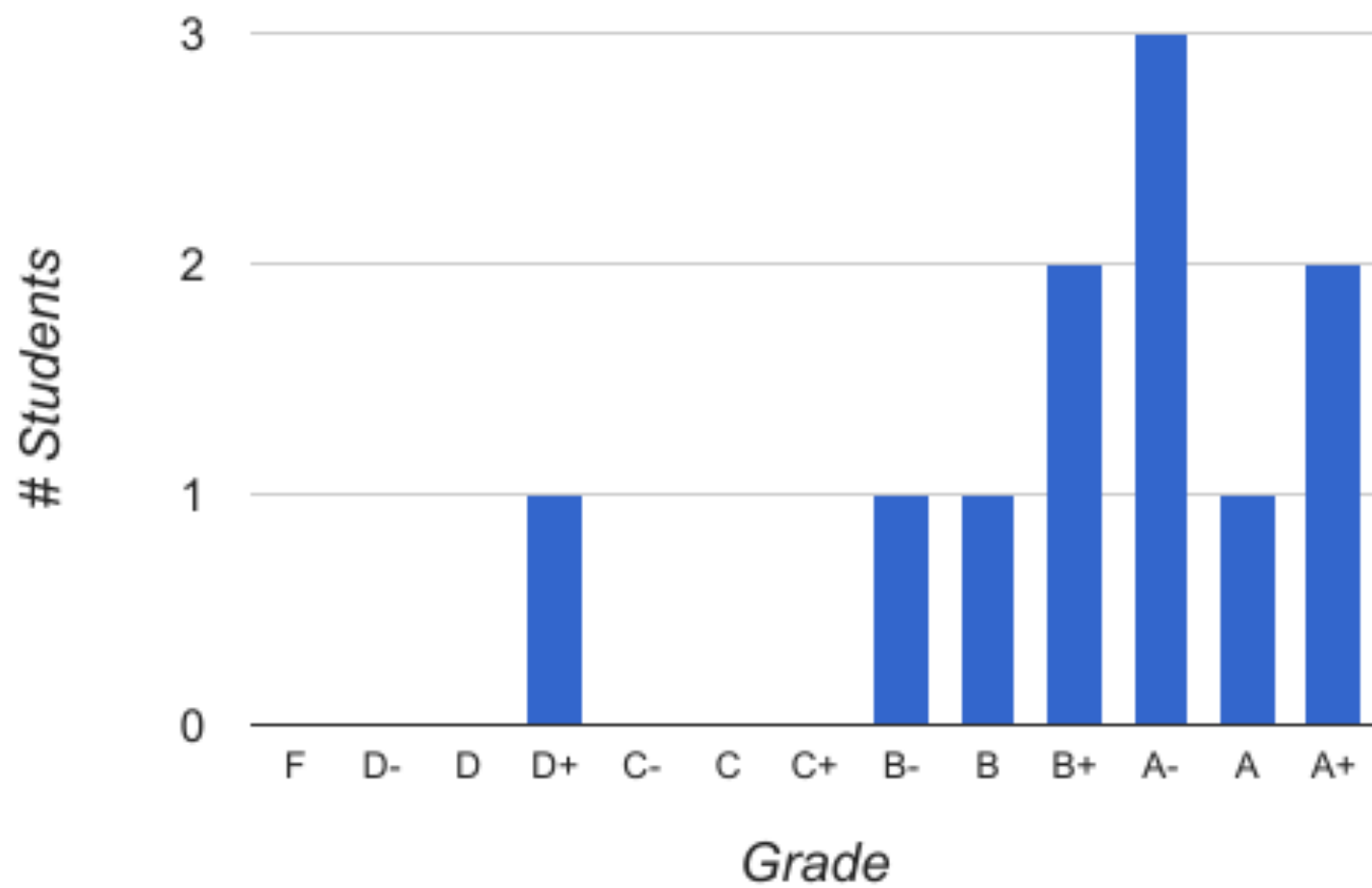
- Self organise groups of 2 for Project 3.

- Project 4 will likely be individual.

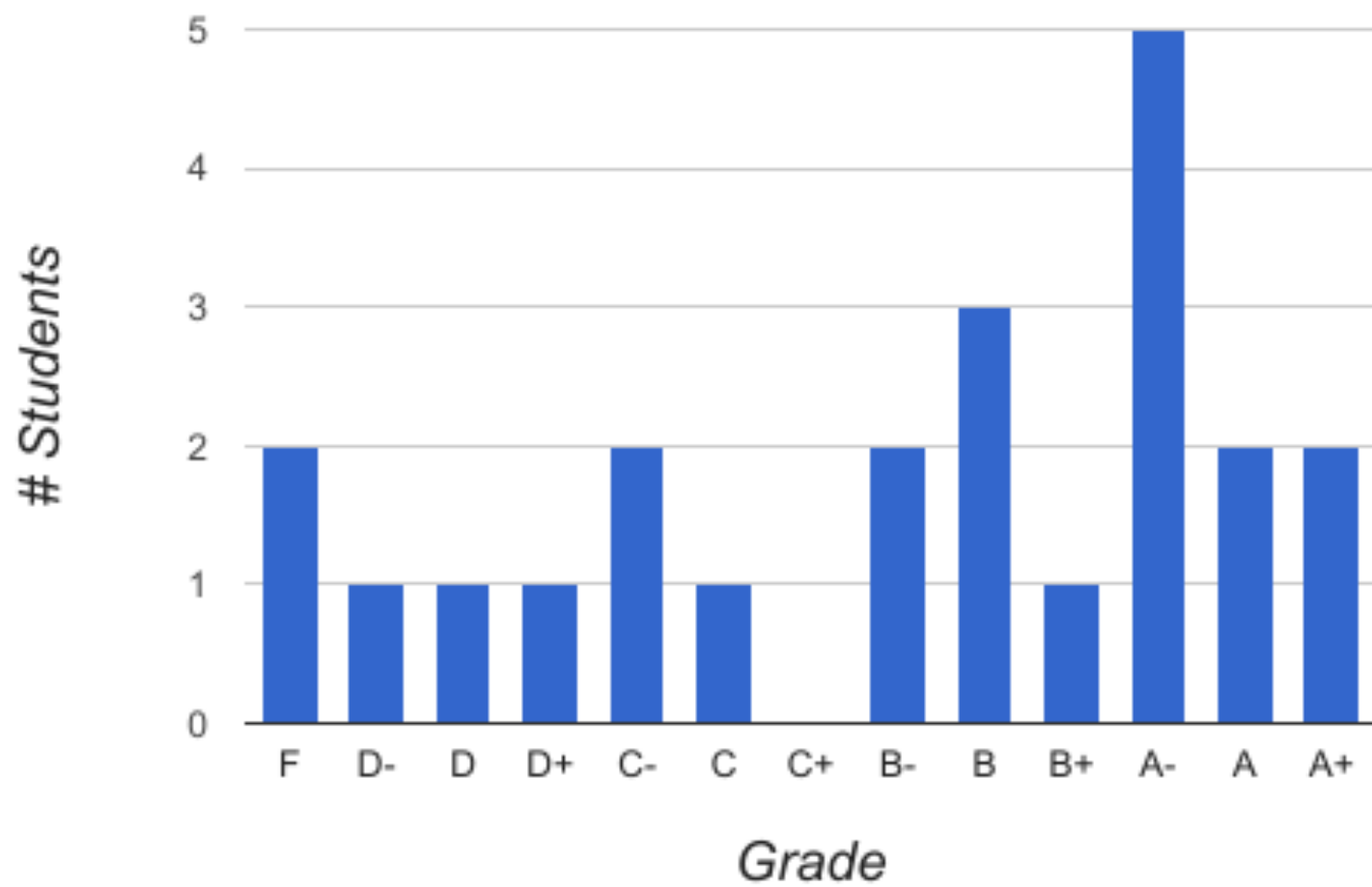
- Transitioning to Modelling and Game Theory this week.



CS423 Midterm Exam Grade Distribution



CS523 Midterm Exam Grade Distribution





Core Wars Class Competition

matthew@electra: /home/matthew/ownCloud/Teaching/CS523/Projects/Project2/Tournament/Entrants

matthew@electra>



matthew@electra>



matthew@electra: /home/matthew/ownCloud/Teaching/CS523/Projects/Project2/Tournament/Entrants

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matthew@electra> █
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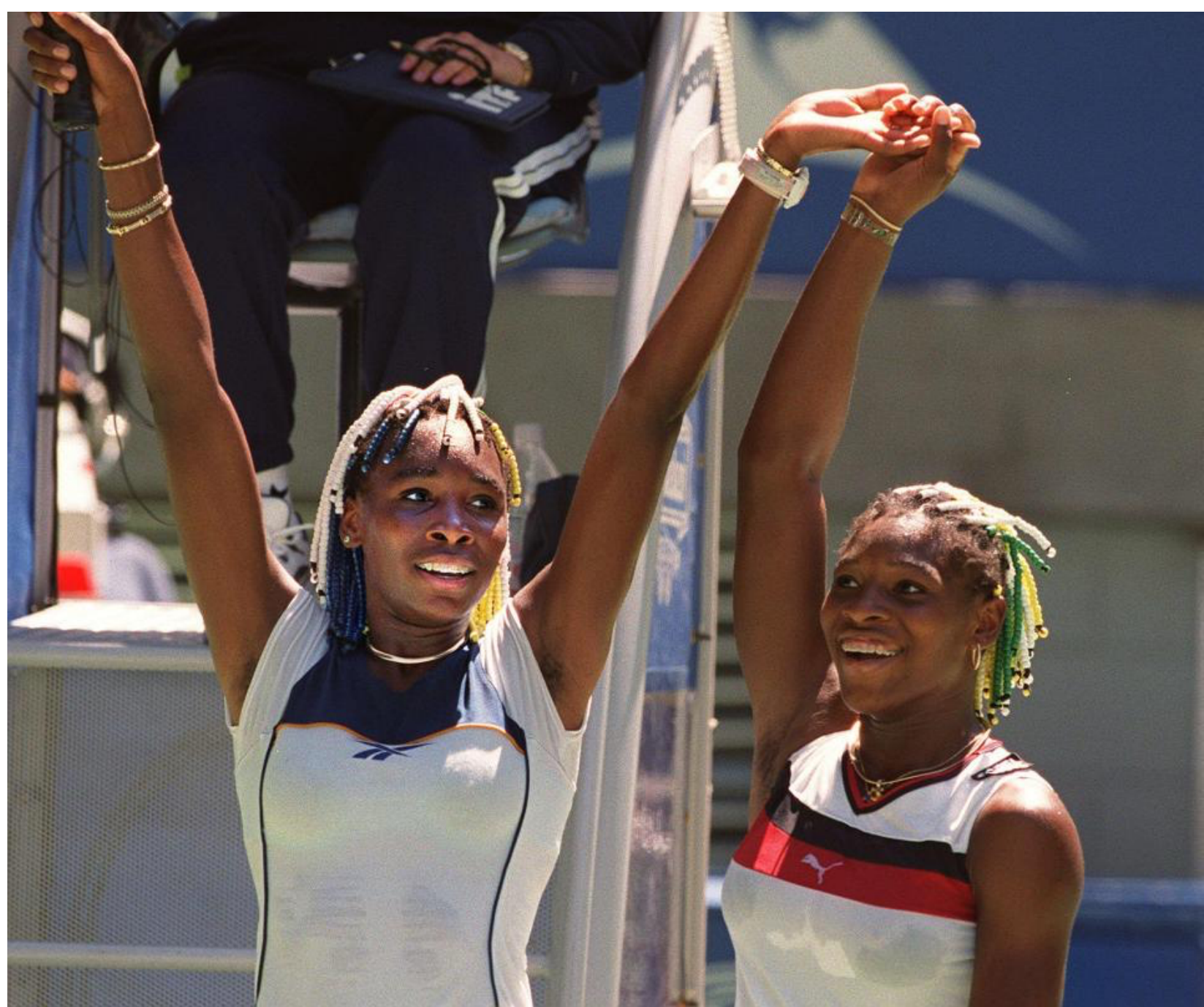


Semifinal: Team 7 vs Team 14 A

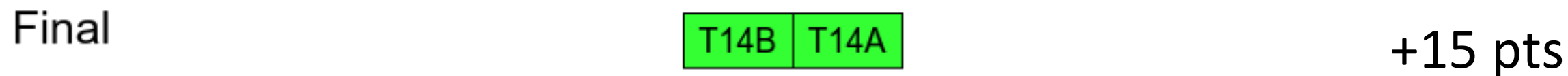


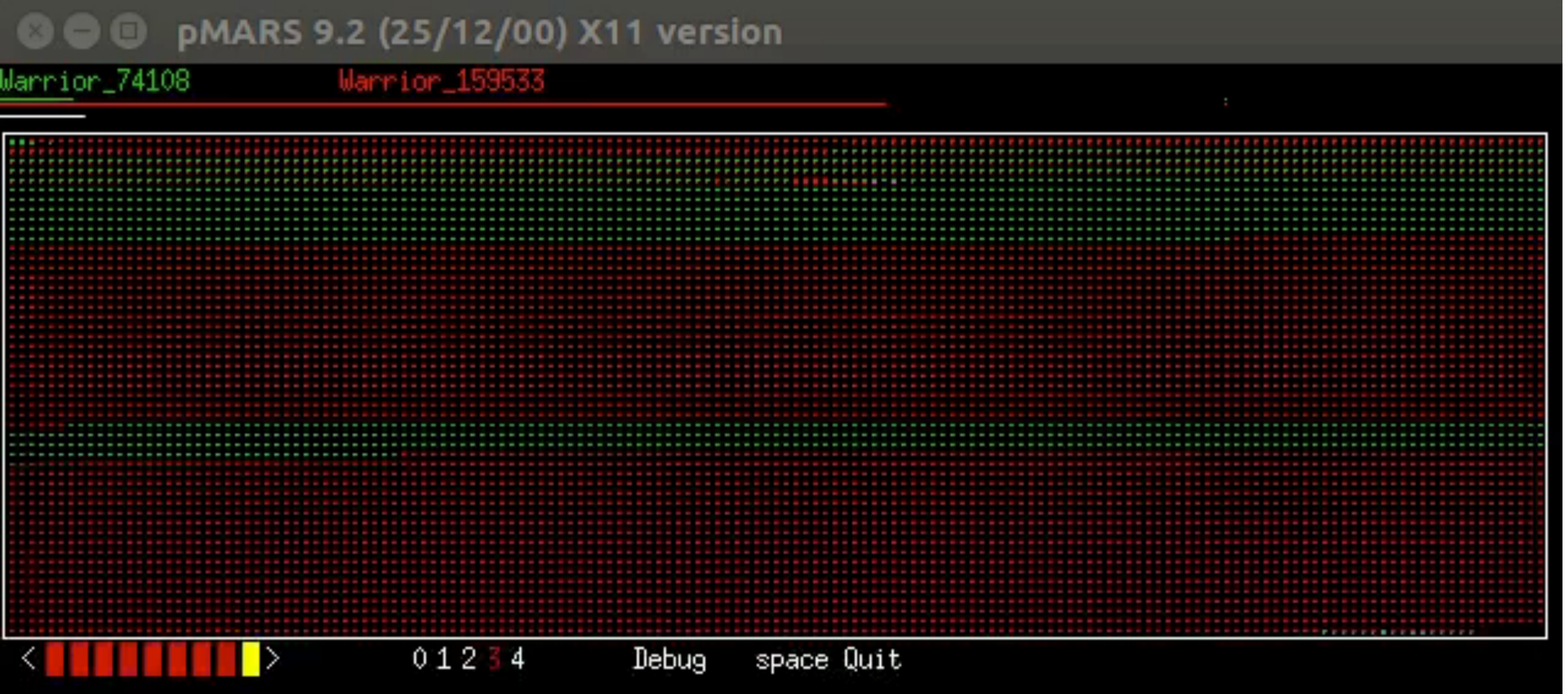
Semifinal: Team 5 vs Team 14 B

Warriors 14 A and B
move to the final!

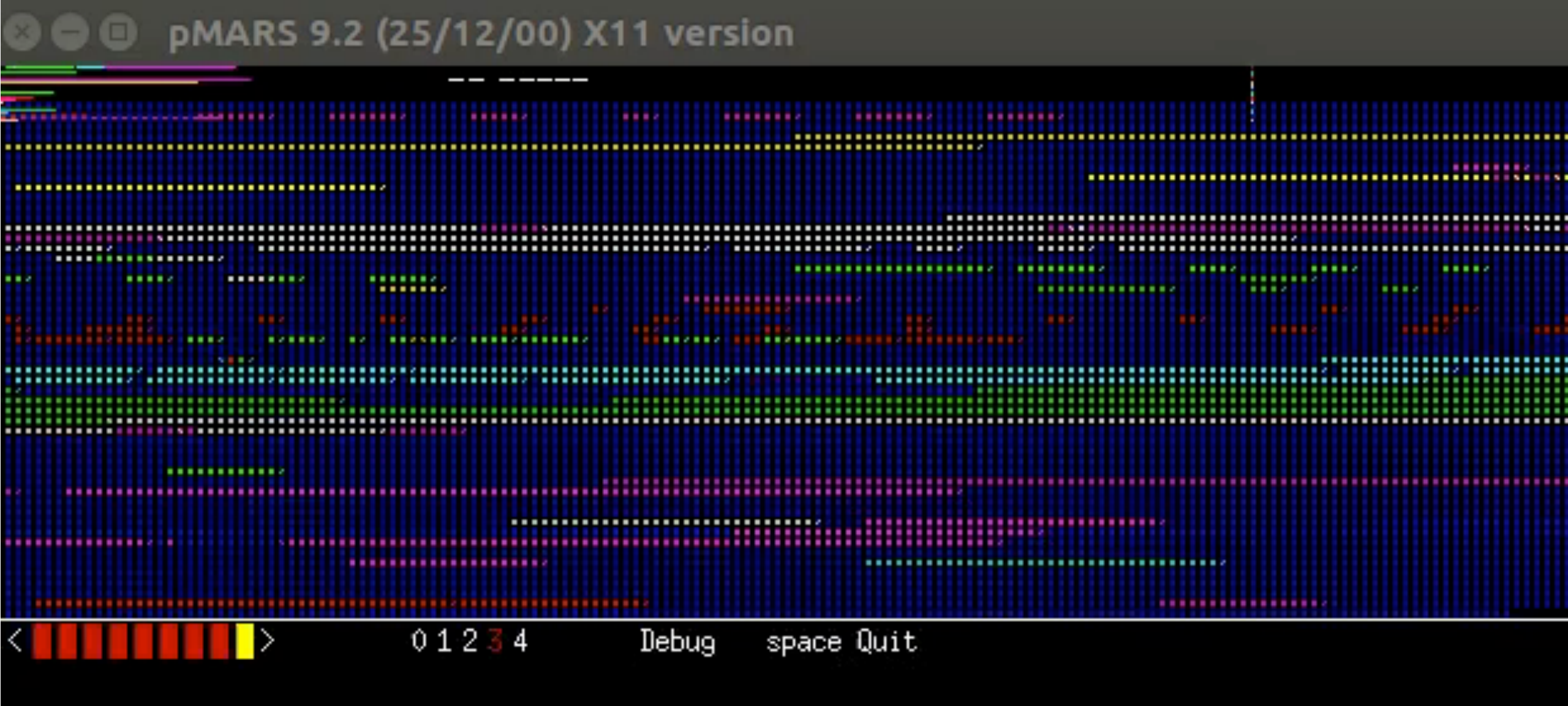


Serena (16) and Venus Williams (17), Australian Open 1st and 2nd place.





Final: Warrior 14A vs 14B (+35 bonus points)



Bonus Round: Free for All

MOV.I * 8, * 5
SUB.AB # 8, * 7

Recent
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Entrants
T01
T02
T03

+10 pts

b21911d9-4ae3-396a-9d54-a2c9b1370424 by T17 scores 49624

Results: 0 0 0 0 0 0 0 0 1 2 3 2 6 9 5 21 24 28 38 51 74 105 129 105 112 76 97 55 24 6 0 0 0 0 28

: T10-blu-solly by : G10 scores 49287

Results: 0 0 0 0 1 0 0 1 0 2 3 3 4 7 2 20 25 26 36 52 73 104 131 106 113 76 97 55 24 6 0 0 0 0 34

: T10-red-solly by : G10 scores 49267

Results: 0 0 0 0 1 0 0 1 0 1 4 3 3 7 3 19 26 28 38 55 72 100 128 106 113 76 97 55 24 6 0 0 0 0 35

DD_203.red by Damion Terrell and Dominic Paul Delvecchio scores 49056

Results: 0 0 0 0 1 0 0 0 0 1 2 3 5 7 5 19 27 28 39 50 71 102 130 105 112 76 95 55 24 6 0 0 0 0 38

c251438d-6d43-36fd-8ff2-0fb8bd5c9f4e by T17 scores 48970

Results: 0 0 0 0 0 0 0 1 0 0 3 2 6 7 4 19 24 29 37 53 71 104 128 106 113 76 97 55 24 6 0 0 0 0 36

Unknown by Anonymous scores 48669

Results: 0 0 0 0 0 0 0 1 0 2 4 3 3 9 4 20 23 25 37 53 72 102 126 103 113 74 97 55 24 6 0 0 0 0 45

DD_204.red by Damion Terrell and Dominic Paul Delvecchio scores 47865

Results: 0 0 0 0 1 0 0 1 0 0 4 1 5 6 4 16 20 25 37 52 71 102 128 104 111 75 96 55 24 6 0 0 0 0 57

Warrior_159533 by T14 scores 47682

Results: 0 0 0 0 0 0 0 1 1 2 0 2 4 7 3 18 22 28 36 50 72 98 128 105 109 74 97 55 24 6 0 0 0 0 59

Warrior_74108 by T14 scores 47518

Results: 0 0 0 0 0 0 0 0 0 2 1 3 3 8 4 19 23 24 37 50 69 100 128 103 111 74 96 55 24 6 0 0 0 0 61

Unknown by Anonymous scores 47495

Results: 0 0 0 0 0 0 0 0 0 2 2 2 4 7 4 19 24 22 35 50 70 99 129 105 109 75 97 55 24 6 0 0 0 0 61

: CURRENT_WARRIOR_NAME by : GROUP_T9 scores 47045

Results: 0 0 0 0 0 0 0 0 0 1 2 2 5 7 1 16 19 30 35 50 69 100 126 104 111 74 97 55 24 6 0 0 0 0 67

: CURRENT_WARRIOR_NAME by : GROUP_T9 scores 46597

Results: 0 0 0 0 0 0 0 0 1 1 0 3 5 4 15 22 28 31 53 67 98 126 106 110 76 96 55 24 6 0 0 0 0 73

~999[706] by T1 scores 46263

Results: 0 0 0 0 0 0 0 0 1 0 3 1 2 6 2 18 18 27 37 50 69 98 124 98 109 75 97 55 24 6 0 0 0 0 81

~999[424] by T1 scores 46243

Results: 0 0 0 0 0 0 0 0 1 0 2 2 3 6 3 12 21 26 35 51 68 95 126 104 110 75 97 54 24 6 0 0 0 0 80

Bester_Neo.RED by Butterfly Wings scores 45286

Results: 0 0 0 0 0 0 0 0 0 1 2 2 3 4 4 15 22 23 32 43 68 101 117 103 112 71 96 54 24 6 0 0 0 0 98

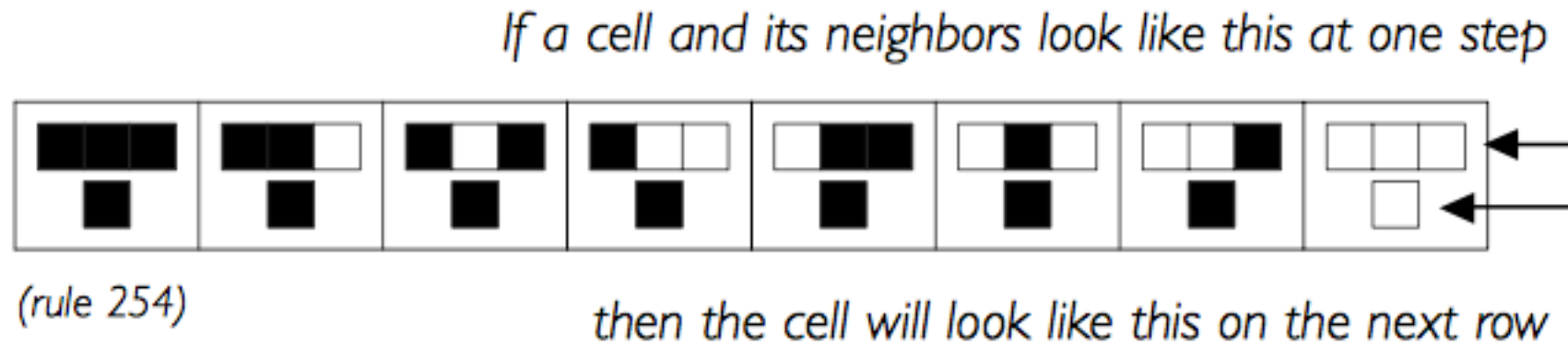
Best_Neo.RED by Butterfly Wings scores 44702

Results: 0 0 0 0 0 0 0 0 0 1 1 2 3 4 4 13 20 22 35 43 64 94 126 101 109 75 94 53 24 6 0 0 0 0 107

: T3-1st-place-warrior.RED by : T3 scores 43857

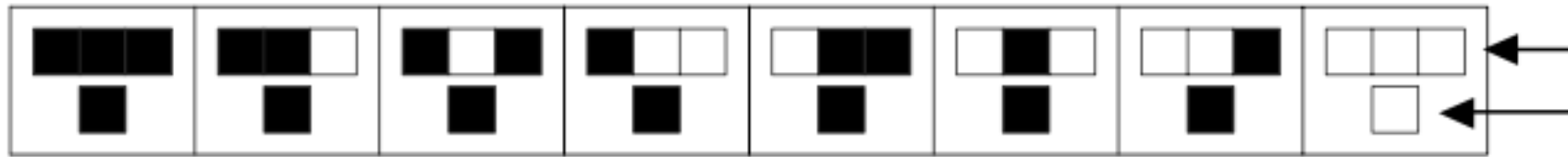
Results: 0 0 0 0 1 0 0 0 0 1 0 1 3 2 17 16 20 24 45 67 98 120 104 109 74 95 55 24 6 0 0 0 0 119

Cellular Automata



Cellular Automata

If a cell and its neighbors look like this at one step



(rule 254)

then the cell will look like this on the next row

With the rule above, a simple pattern is produced.



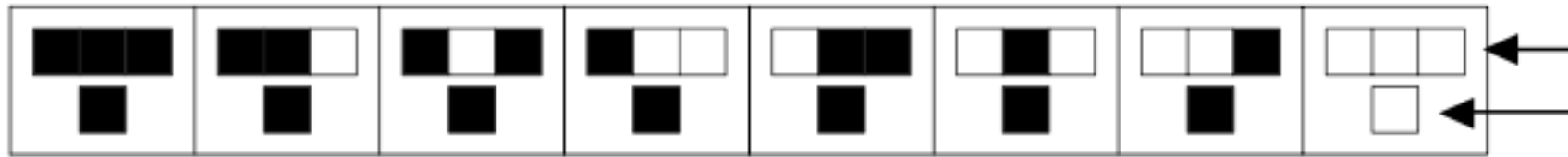
step 1



step 2

Cellular Automata

If a cell and its neighbors look like this at one step



(rule 254)

then the cell will look like this on the next row

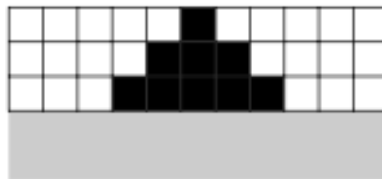
With the rule above, a simple pattern is produced.



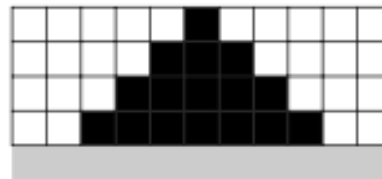
step 1



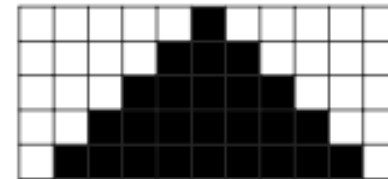
step 2



step 3



step 4



step 5

Cellular Automata

Wolfram's Rule Numbering System.

We keep the order of input combinations fixed.

Our numbering is just the decimal for the output bits:

111

110

101

100

011

010

001

000

Cellular Automata

Wolfram's Rule Numbering System.

We keep the order of input combinations fixed.

Our numbering is just the decimal for the output bits:

111	?
110	?
101	?
100	?
011	→ ?
010	?
001	?
000	?

Cellular Automata

Wolfram's Rule Numbering System.

We keep the order of input combinations fixed.

Our numbering is just the decimal for the output bits:

111	1	
110	1	
101	1	
100	1	
011	1	\rightarrow
010	1	
001	1	
000	0	

$= 11111110 = 254$

Cellular Automata

Wolfram's Rule Numbering System.

We keep the order of input combinations fixed.

Our numbering is just the decimal for the output bits:

111	0	
110	0	
101	0	
100	1	$\rightarrow 1 = 00011110 = 30$
011	1	
010	1	
001	1	
000	0	

Cellular Automata

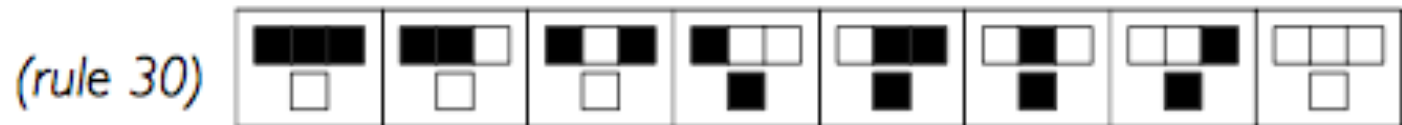
Wolfram's Rule Numbering System.

We keep the order of input combinations fixed.

Our numbering is just the decimal for the output bits:

111	0
110	0
101	0
100	1
011	1
010	1
001	1
000	0

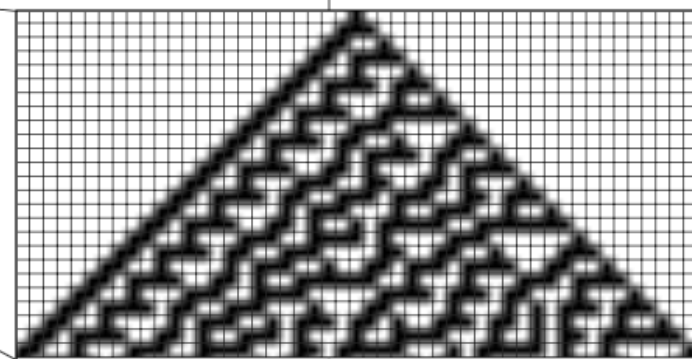
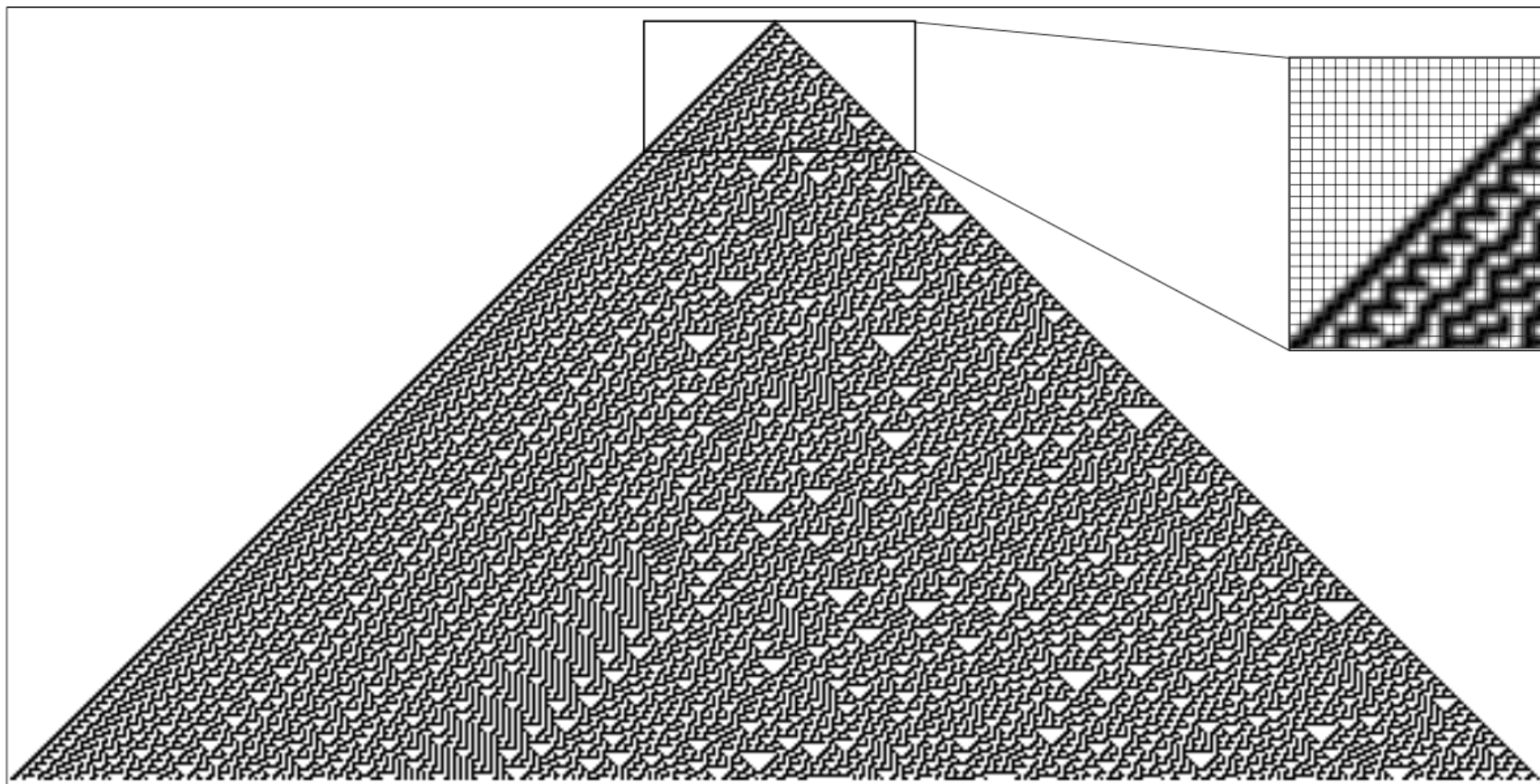
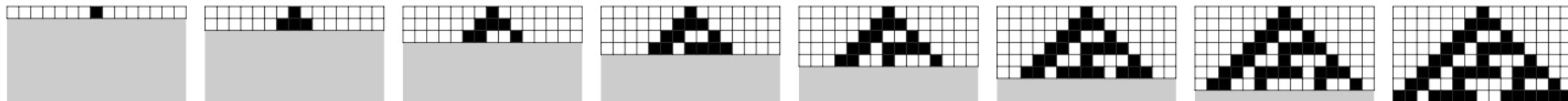
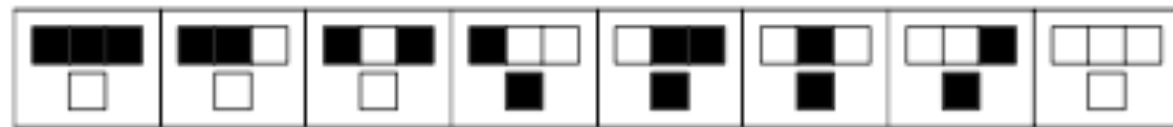
$\rightarrow 1 = 00011110 = 30$



Number of Elementary CAs

- We have seen that there are 256 unique rules.
- We can reduce this number to 128 because of the 1/0 symmetry.
- We can further reduce this number to 64 through left/right symmetry.

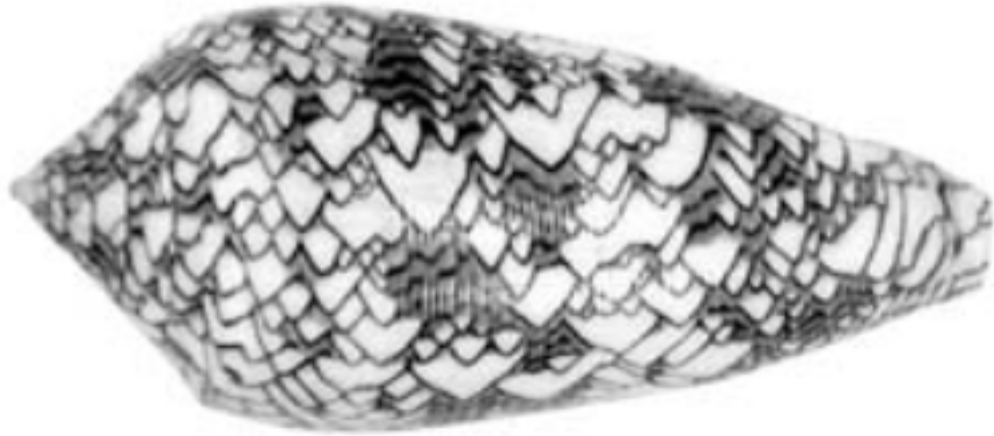
(rule 30)



25 steps

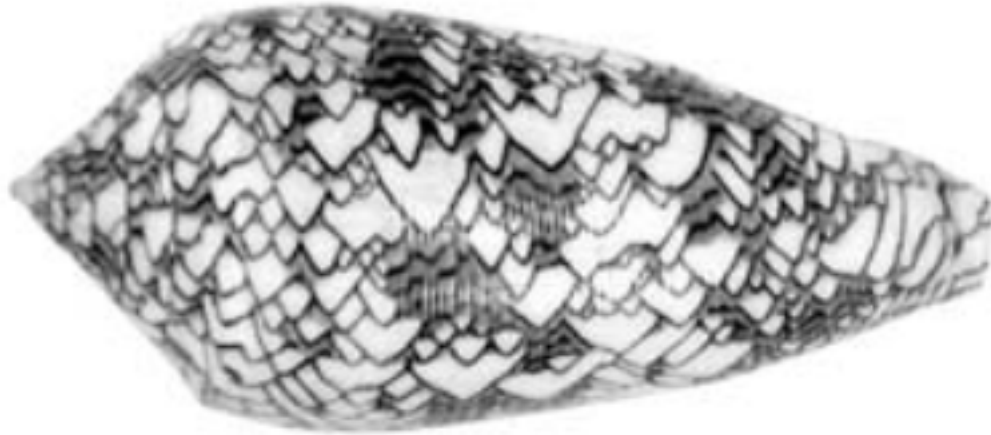
250 steps

Cellular Automata



Seashell pattern

Cellular Automata



Seashell pattern

We are back to our recurring theme in this class. Simple rules can produce very complex behaviour. The complexity of the world (computational and physical) can be modelled with simple rules, and that This may be the **only** way to study many systems.

Universal Computation

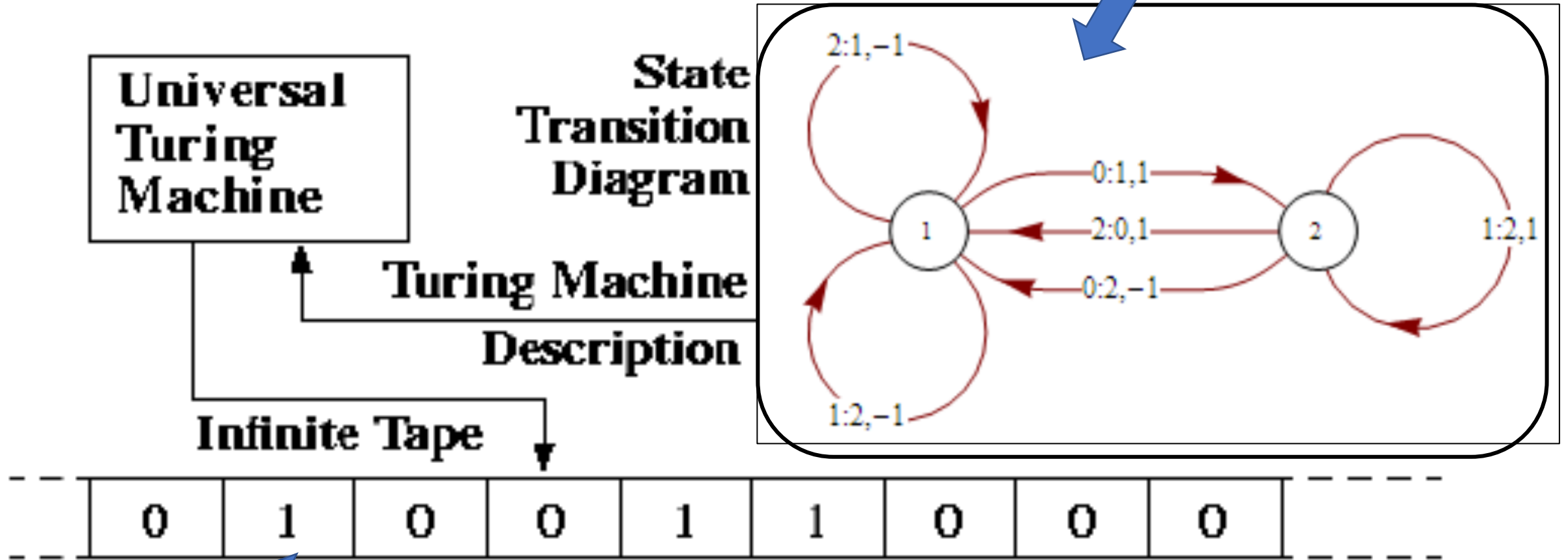
- One of the most powerful systems we know of is the Universal Turing Machine (The subject of most of CS500).
- Each Turing Machine is capable of solving an problem. Modelled after an abstract mathematician (who has paper, pencil, and lookup rules).
- Allan Turing showed that there is a single Turing Machine that can simulate every other Turing Machine.
- So from simple lookup rules, pencil and paper we can generate the entire diversity of the computational universe.
- Von Neumann's self-reproducing automata was the first universal CA.

What is a program really?

- A set of rules
- An input state
- An output

Universal Turing Machine

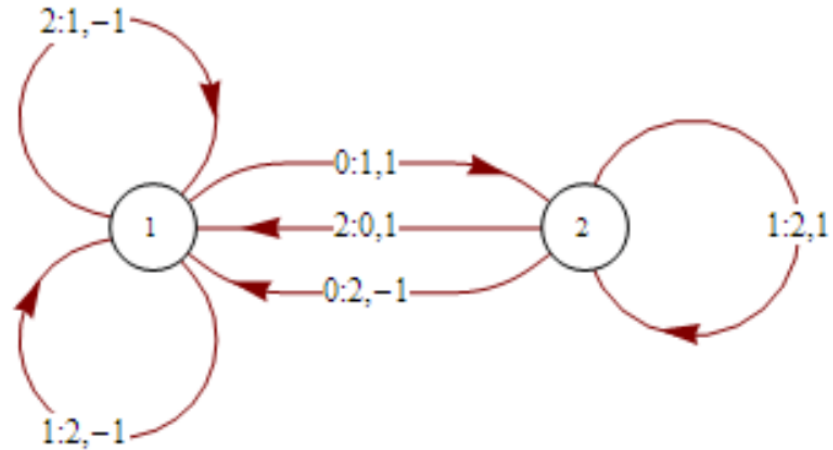
Rules



Initial value of the tape is the input

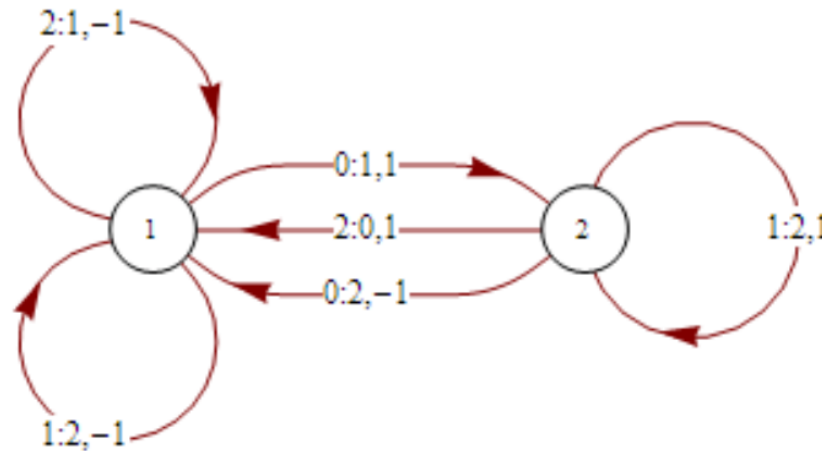
Universal Computation

- The “brain” of the mathematician in a Turing Machine is a finite state automata:



Universal Computation

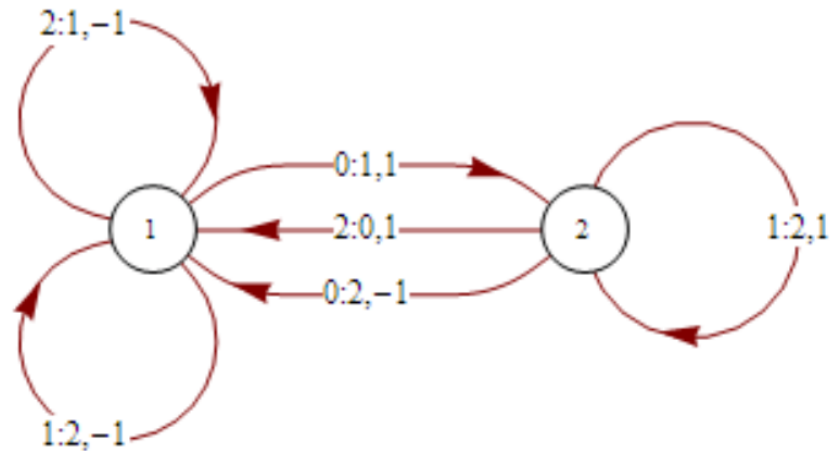
- The “brain” of the mathematician in a Turing Machine is a finite state automata (FSA):



- Wolfram showed that you can encode these FSA as rules in a cellular automata.

Universal Computation

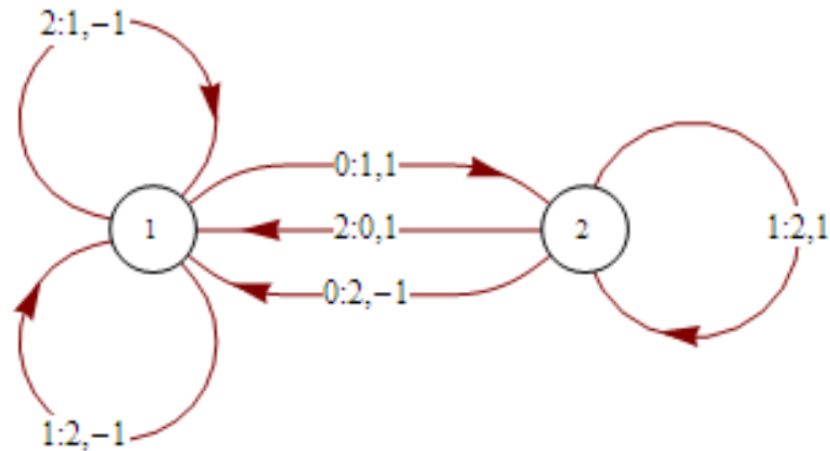
- We are no longer thinking about elementary CAs where we just look at the neighbouring cells.



- Wolfram showed that you can encode these FSA as rules in a cellular automata.

Universal Computation

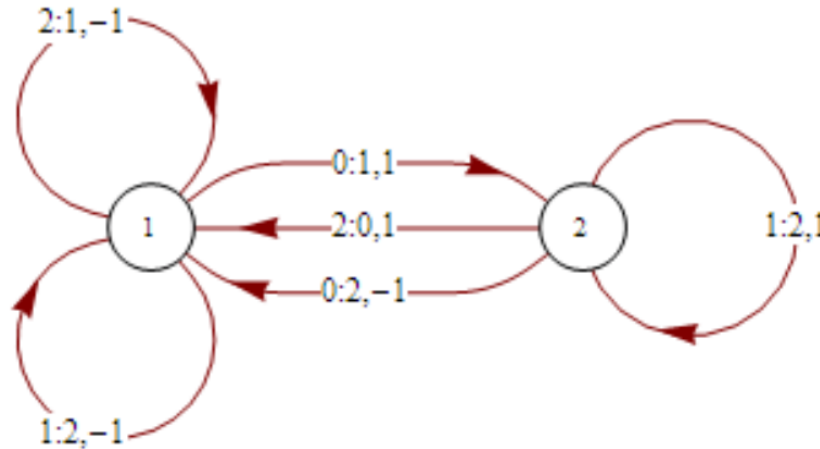
- We are no longer thinking about elementary CAs where we just look at the neighbouring cells.



- In this example there are only 3 symbols and 2 states in the FSA. The 1 and -1 after the comma indicate the offset.

Universal Computation

- For elementary machines we saw there are 2^8 rules for 2^3 inputs.

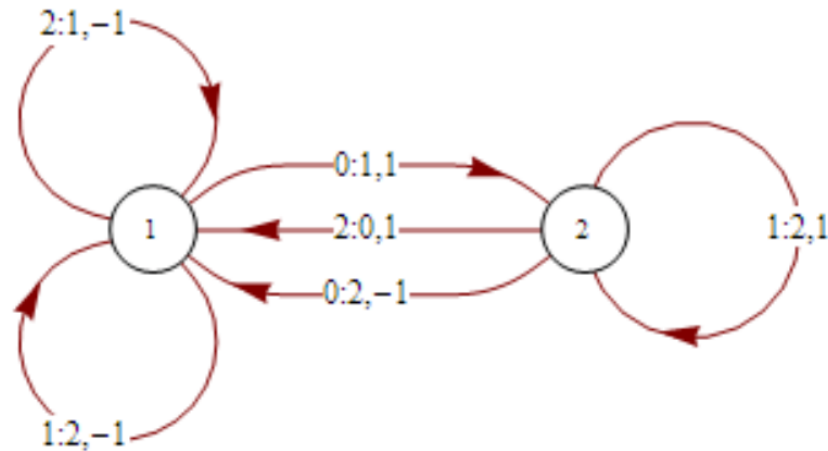


- The number of Wolfram's Turing Machine CAs is larger

$$(2|\Sigma||S|)^{2 \cdot 3} = 12^6 = 2985984$$

Universal Computation

- This is machine number 596440 in Wolfram's scheme.



- The number of Wolfram's Turing Machine CAs is larger

$$(2|\Sigma||S|)^{|\Sigma||S|} = 12^6 = 2985984$$

Translating a Turing Machine into a CA.

	In State A	In State B
0	Print 1, Move Right, Goto State B	Print 2, Move Left, Goto State A
1	Print 2, Move Left, Goto State A	Print 2, Move Right, Goto State B
2	Print 1, Move Left, Goto State A	Print 0, Move Right, Goto State A

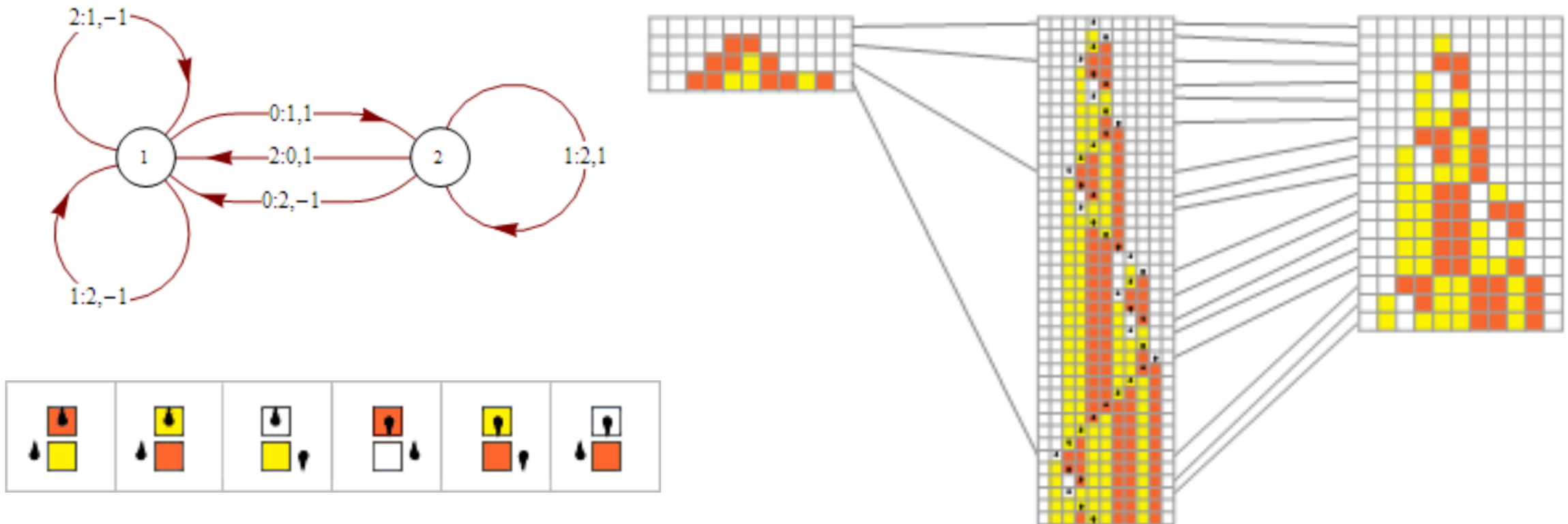


Colours = symbols on the tape
 Each column is the transition from tape input to output
 The droplet indicates whether the tape read head moved left or right.

Claimed but not proven that this is the smallest possible universal Turing machine (the proof by Alex Smith is being debated)

Universal Computation

- This is machine number 596440 in Wolfram's scheme.



Cellular Automata

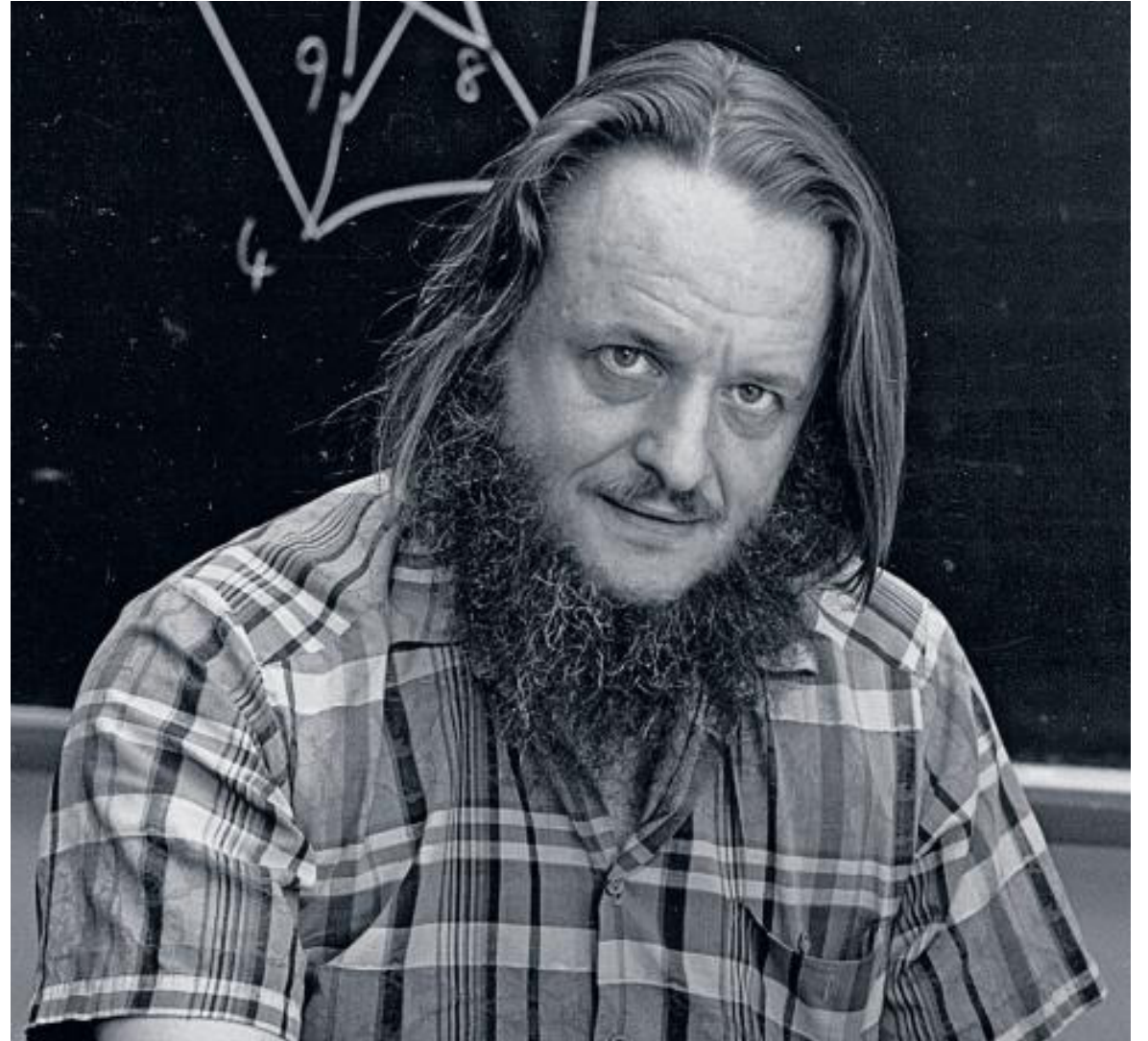
One-dimensional cellular automata can be universal.

Wolfram (New Kind of Science, 2002) gave an example of a 19-color universal one-dimensional next-nearest neighbor cellular automaton in which a block of 20 cells is used to represent each single cell in the cellular automaton being emulated.

Even more amazing [rule 110 elementary cellular automaton](#) is universal. (Cook, M. "Universality in Elementary Cellular Automata." *Complex Systems* **15**, 1-40, 2004.).

2D Automata: Game of Life

- John Conway was interested in simplifying John von Neumann's rules for self-reproducing automata.
- Invented a 2D cellular automata to demonstrate this in 1970.



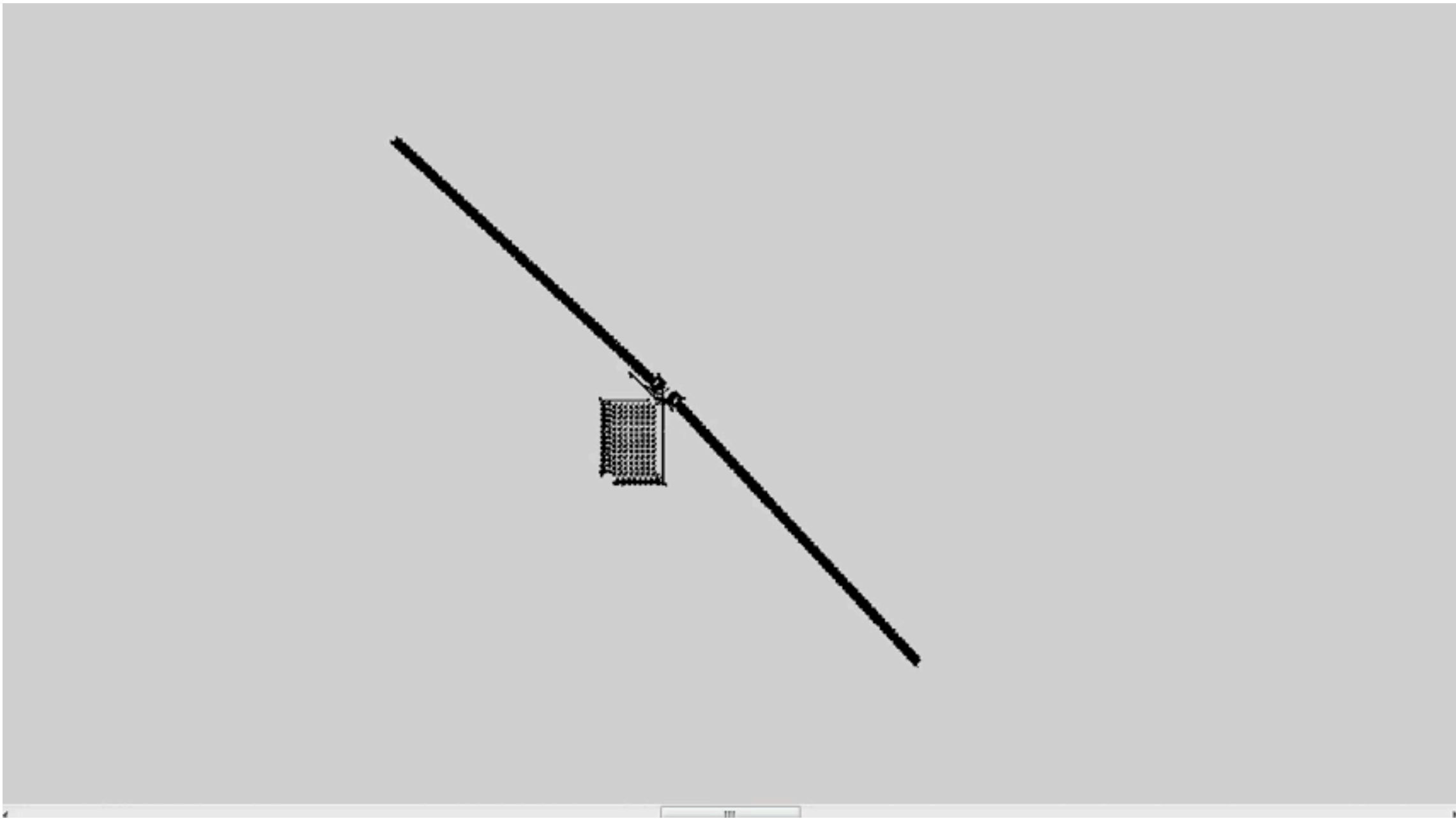


Numberphile: <https://www.youtube.com/watch?v=R9Plq-D1gEk>

2D Automata: Game of Life - Rules

- Any live cell with fewer than 2 live neighbours dies
- Any live cell with more than 3 live neighbours dies
- Any dead cell with exactly 3 live neighbours becomes alive
- Otherwise (2 or 3 neighbors) the cell is unchanged

Alex Bellos: <https://www.youtube.com/watch?v=vGWGeund3eA>



Models

- We have seen a lot of different models in this class being used for different things.
- But what is a model? What can they do?

Models

Can you think of a model we have seen in class that address each of the following. (Work with the person next to you)

- A model that makes *predictions* about some system
- A models used to define *computation*
- *Existence proof* models (models demonstrating the possibility of something).
- A model used to *explain* something that already happened.



Models

- `Now it would be very remarkable if any system existing in the real world could be *exactly* represented by any simple model. However, cunningly chosen parsimonious models often do provide remarkably useful approximations.`
- `For such a model there is no need to ask the question "Is the model true?". If "truth" is to be the "whole truth" the answer must be "No". The only question of interest is "Is the model illuminating and useful?"`.

Box, G. E. P. (1979), "Robustness in the strategy of scientific model building", in Launer, R. L.; Wilkinson, G. N., Robustness in Statistics, Academic Press, pp. 201–236.

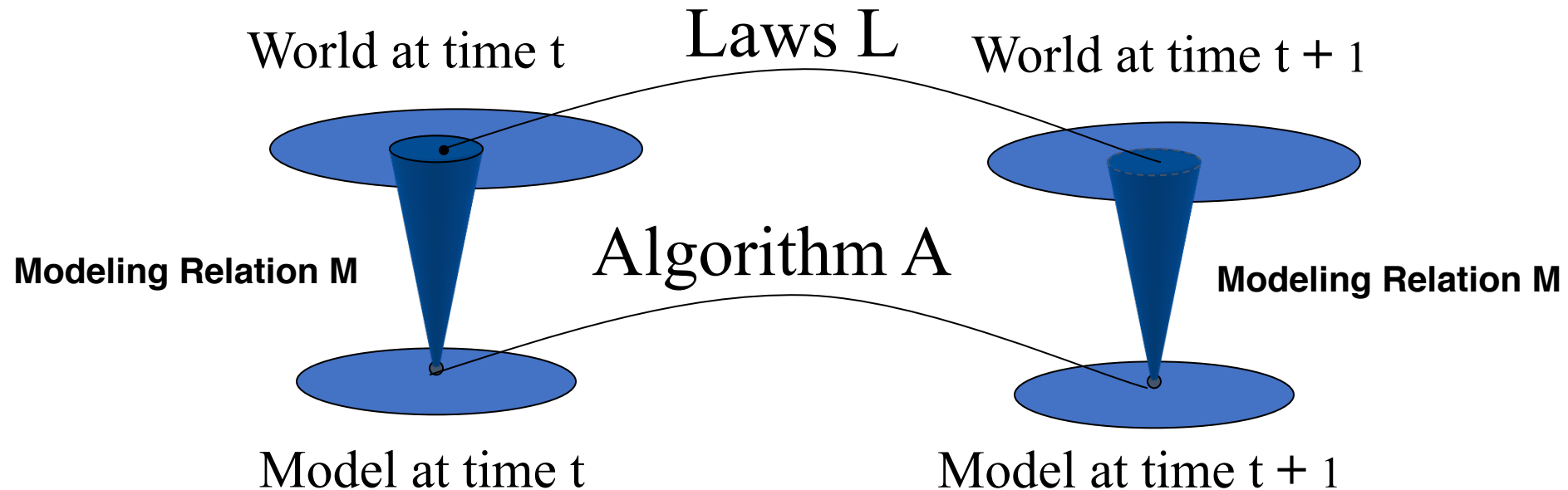
Models

- `Now it would be very remarkable if any system existing in the real world could be *exactly* represented by any simple model. However, cunningly chosen parsimonious models often do provide remarkably useful approximations.`
- `For such a model there is no need to ask the question "Is the model true?". If "truth" is to be the "whole truth" the answer must be "No". The only question of interest is "Is the model illuminating and useful?"`.
- “All models are wrong, some are useful.”

Box, George. E. P. (1979), "Robustness in the strategy of scientific model building", in Launer, R. L.; Wilkinson, G. N., Robustness in Statistics, Academic Press, pp. 201–236.

Models as Homomorphic Maps

Commutativity of the Diagram



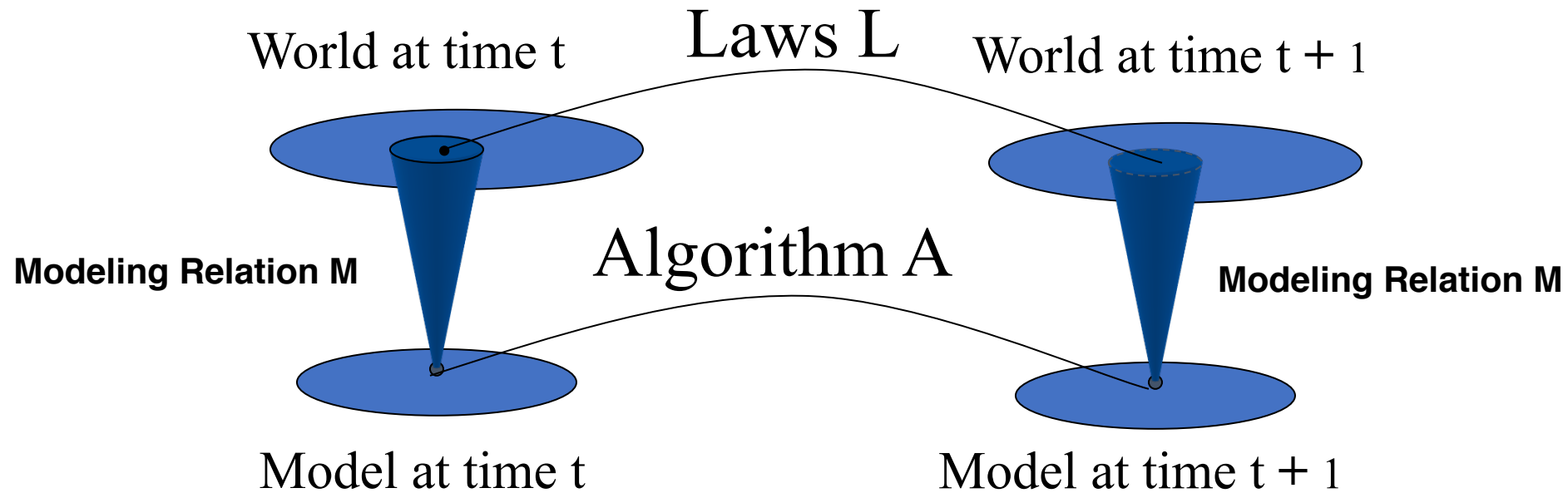
M is an equivalence relation.

Model M is valid if this is a homomorphic map:

$$M(L(x)) = A(M(x))$$

Models as Homomorphic Maps

transformation of one set into another that preserves in the second set the relations between elements of the first.



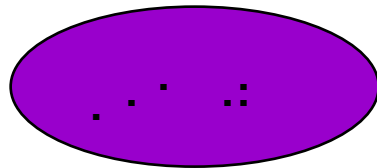
Models

- “It can scarcely be denied that the supreme goal of all theory is to make the irreducible basic elements as simple and as few as possible without having to surrender the adequate representation of a single datum of experience.”

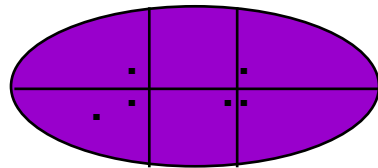
Attributed to Albert Einstein in “On the Method of Theoretical Physics,” the Herbert Spencer Lecture, Oxford, June 10, 1933. This is the Oxford University’ Press

Equivalence Classes (CS261)

- Equivalence class =
 $\{x \mid x \in R\}$ and R is an equivalence relation.
- R is an equivalence relation:
 - Reflexive: $\forall x(xRx)$
 - Symmetric: $xRy \Rightarrow yRx$
 - Transitive: $(xRy) \wedge (yRz) \Rightarrow (xRz)$
- Example: $xRy \Leftrightarrow x$ and y are in the same little box.



Set of Objects



Partition set into 6 little boxes



Equivalence classes

Examples of Equivalence Relations

- “Is similar to” or “congruent to” on the set of all triangles.
- Logical equivalence of statements in logic.
- “Has the same image under a function” on the elements of the domain of the function.
- What’s not an equivalence relation?
 - The relation “ \geq ” between real numbers is reflexive and transitive, but not symmetric. For example, $7 \geq 5$ does not imply that $5 \geq 7$. It is, however, a partial order.
 - The relation “is a sibling of” on the set of all human beings is not an equivalence relation.
 - Is Symmetric (if A is a sibling of B, then B is a sibling of A)
 - Not reflexive (no one is a sibling of himself),
 - Not transitive (since if A is a sibling of B, then B is a sibling of A, but A is not a sibling of A).

Example Homomorphism:

Multiplication of Integers

- Model all pairs of integers and their product:
 - e.g., $14792 \times 4183584 = 61883574528$
- Model:
 - Even X Even = Even
 - Even X Odd = Even
 - Odd X Even = Even
 - Odd X Odd = Odd

Example Homomorphism:

Multiplication of Integers

Model:

Even x Even = Even

Even x Odd = Even

Odd x Odd = Odd

Odd x Even = Even

$$M(L(x)) = M(2^n \times 2^m = 2^k) = \text{Even} \times \text{Even} = \text{Even}$$

The relationships are preserved under our model.

Model relationship:

$$2^n \times 2^m = 2^k \ R \ \text{Even} \times \text{Even} = \text{Even}$$

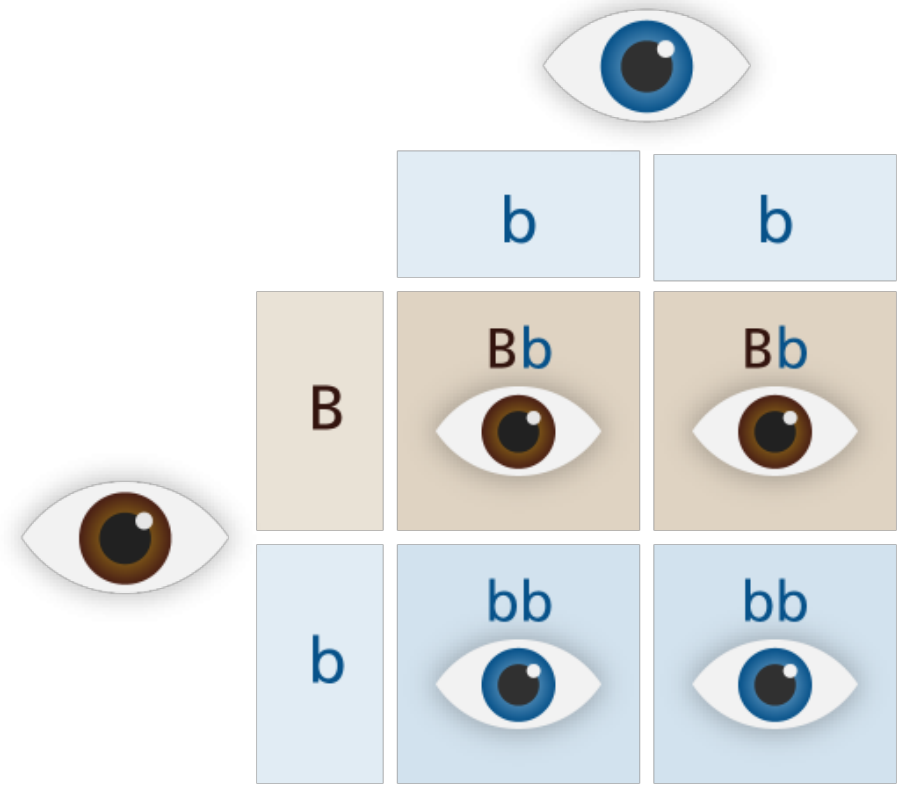
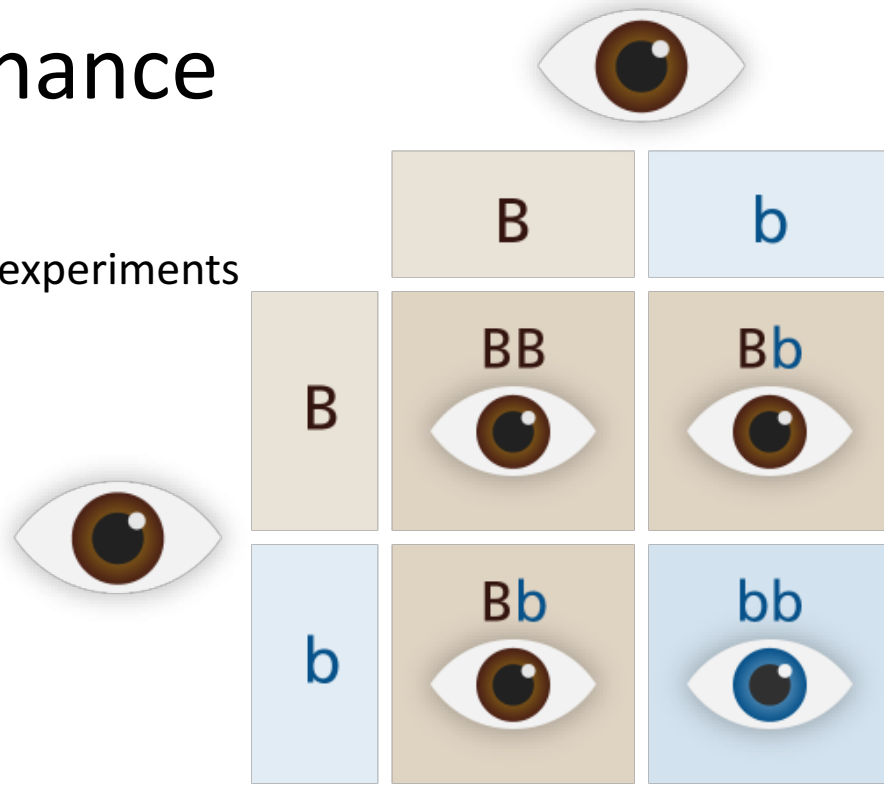
$$2^{n+1} \times 2^{m+1} = 2^{k+1} \ R \ \text{Odd} \times \text{Odd} = \text{Odd}$$

$$2^n \times 2^{m+1} = 2^k \ R \ \text{Even} \times \text{Odd} = \text{Even}$$

$$2^{n+1} \times 2^m = 2^{k+1} \ R \ \text{Odd} \times \text{Even} = \text{Even}$$

Allele Dominance

Recall Gregor Mendel's experiments with peas



B - dominant brown eye allele

b - recessive blue eye allele

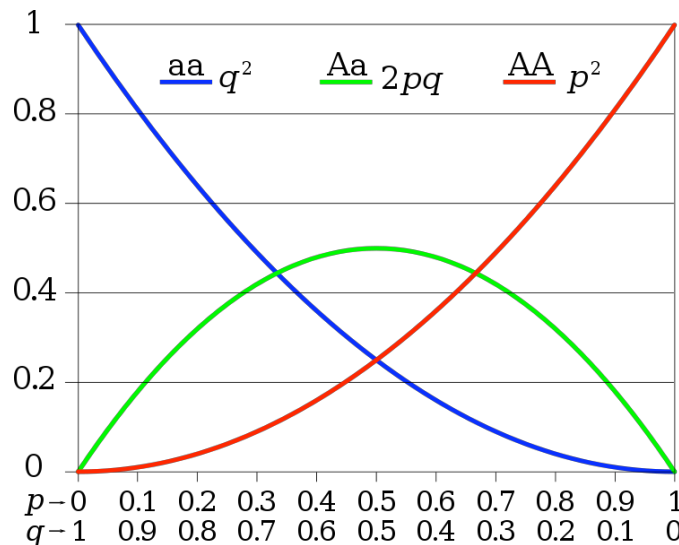
BB brown eyes

Bb brown eyes

bb blue eyes

Example: Hardy-Weinberg Equilibrium

- Allele and genotype frequencies in a population remain constant from generation to generation in the absence of disturbing influences, such as:
 - Non-random mating, mutation, selection, limited population size, overlapping generations, random genetic drift, gene flow, and meiotic drive
- Impossible in nature. Genetic equilibrium is an ideal state that provides a baseline against which to measure change.



$freq(A) = p; freq(a) = q; p + q = 1$. If the population is in equilibrium, then we will have $freq(AA) = p^2$ for the AA homozygotes in the population, $freq(aa) = q^2$ for the aa homozygotes, and $freq(Aa) = 2pq$ for the heterozygotes.

Source: Wikipedia

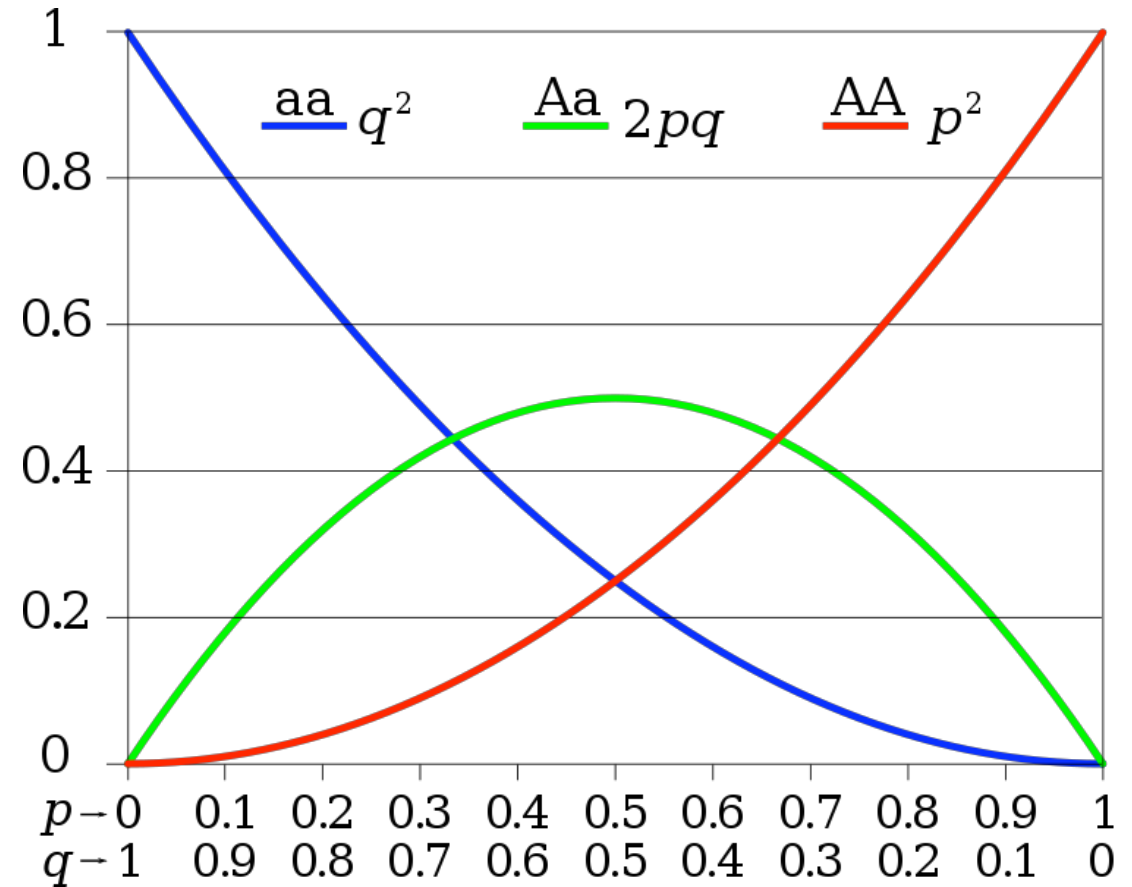
Homomorphic Model (sketch)

Hardy-Weinberg Equilibrium

- Taken from Melanie Moses (2000).
- Example CAS:
 - Theoretical population of desert shrubs, randomly distributed across a homogeneous landscape and isolated from all other populations of shrubs.
 - The shrubs have one relevant genetic variation with two values (alleles).
 - Genotype AA
 - Thrives in annual mean high temperatures T ($29 \leq T \leq 33$).
 - Let X = number of AA individuals.
 - Genotype Aa (and aA)
 - Thrives in annual mean high temperature $28 \leq T \leq 32$.
 - Let Y = number of Aa and aA individuals.
 - Genotype aa
 - Thrives where $27 \leq T \leq 31$.
 - Let Z = number of aa individuals.
 - Need some rules for how populations change (not shown).
- Set of states for system are all (X, Y, Z, T) s.t. $X+Y+Z=N$ for all T :
 - Fixed size population.

Hardy-Weinberg Equilibrium cont.

- Equivalence classes for this set are all X, Y, Z, T such that the ratio of $X:Y:Z$ is constant and T is constant:
 - E.g., $\langle 100, 200, 100, 30 \rangle$ and $\langle 200, 400, 200, 30 \rangle$ are in the same class.
 - Can show that this defines an equivalence relation.
- Transition rules for model are of the form:
 - If $29 \leq T \leq 31$, then population remains in Hardy-Weinberg equilibrium: $X=p^2$, $Y=2pq$, and $Z=q^2$. (p =proportion of A 's in population, $q = 1-p$).
 - If $T=32$: $X=p^2 + 0.05q^2$; $Y = 2pq + 0.05q^2$; $Z = 0.9q^2$
 - Etc.
- Now, show that this is a homomorphism.



Source: Wikipedia

Homomorphic Models Final Remarks

- What if the inside of the box is not deterministic?
 - Search for a set of inputs/outputs that take more variables into account (refine the model) and see if the result is deterministic. OR
 - Look for statistical determinacy, i.e. determinacy in the averages, as in Markov chains.
- Models themselves are seldom regarded in all their detail:
 - Only *some* aspect of the model is related to the system of study.
 - E.g., a tin mouse may be a satisfactory model of a living system, if we ignore the “tinniness” of the model and the “proteinness” of the living mouse.
 - Thus, two systems (the system of study and the model) are often related such that a homomorphism of one is isomorphic with a homomorphism of the other.

Lattice Gas Models (LGCA)

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- That is, we can use continuous values of pressure, temperature, and velocity

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 - We have to model the individual molecules of gas.
- If we are modelling systems with very high energies (such as a nuclear explosion) we have to have a discrete model of the internal states of the atoms involved.

Lattice Gas Models

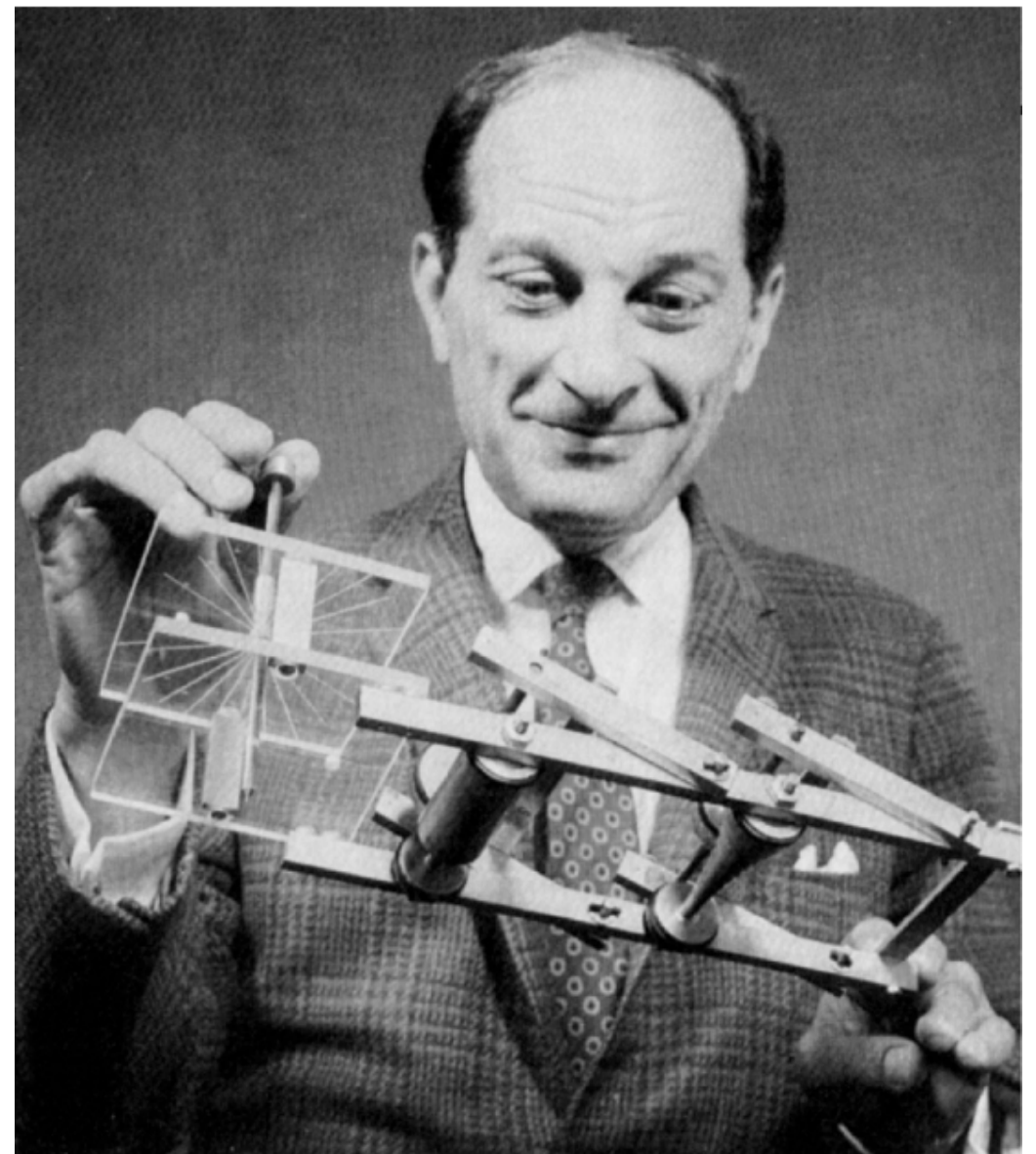
- If the molecules are very cold quantum effects start to dominate their interactions.
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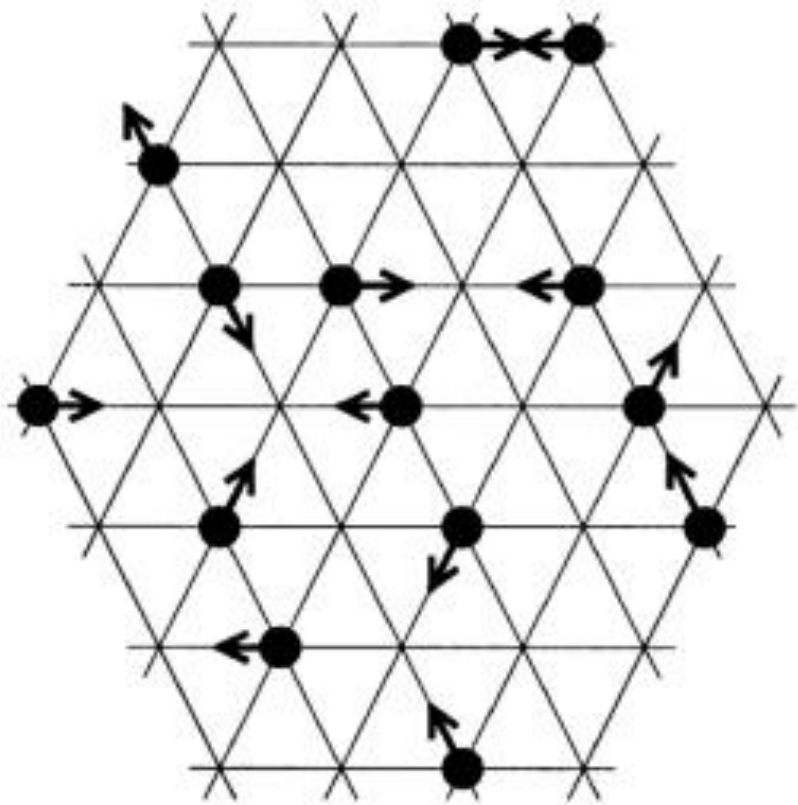
- If the molecules are very cold quantum effects start to dominate their interactions.
 - Here we have to model the quantum effects explicitly.
- These systems require models of the **microscopic** behaviour
- Models that are able to describe the behaviour of the system using just pressure, velocity, and temperature are **macroscopic**.
- Of course, we could model all gasses and fluids at the microscopic level.

Lattice Gas Models

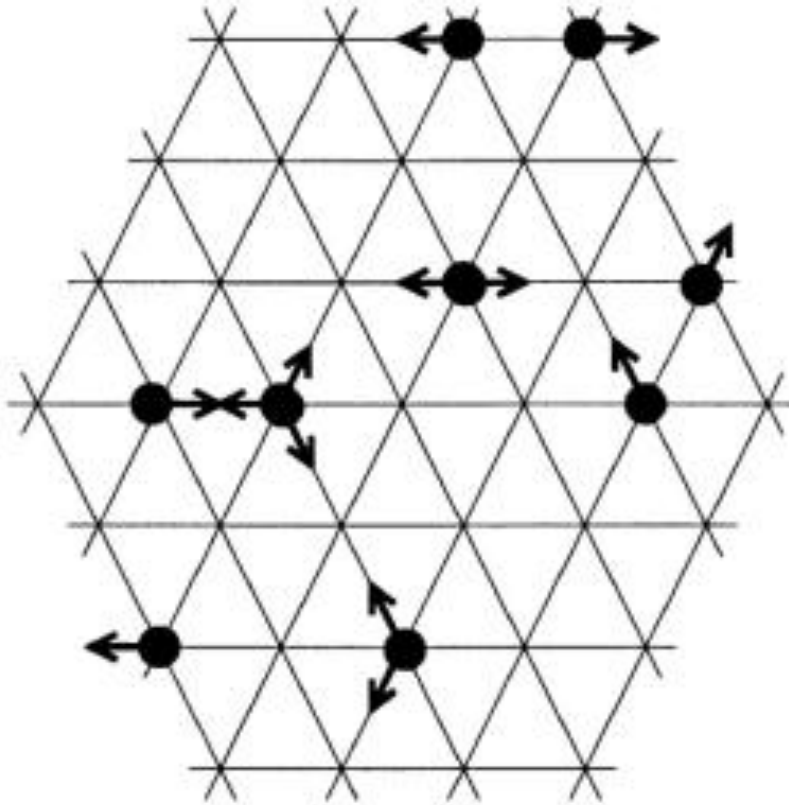
- Cellular automata are used to model molecular systems.
- The use of cellular automata to model particles such as gasses, fluids, and
- The propagation of subatomic particles was pioneered by Stanislaw Ulam and John von Neumann in the 1950s.



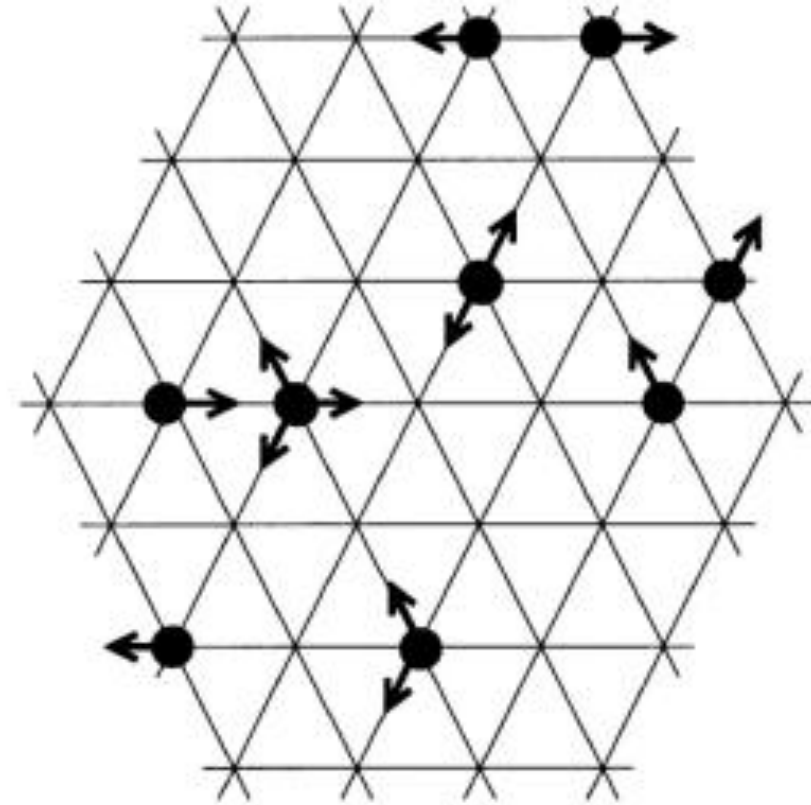
Stanislaw Ulam with the FERMIAC, used to model neutron transport, Los Alamos National Labs



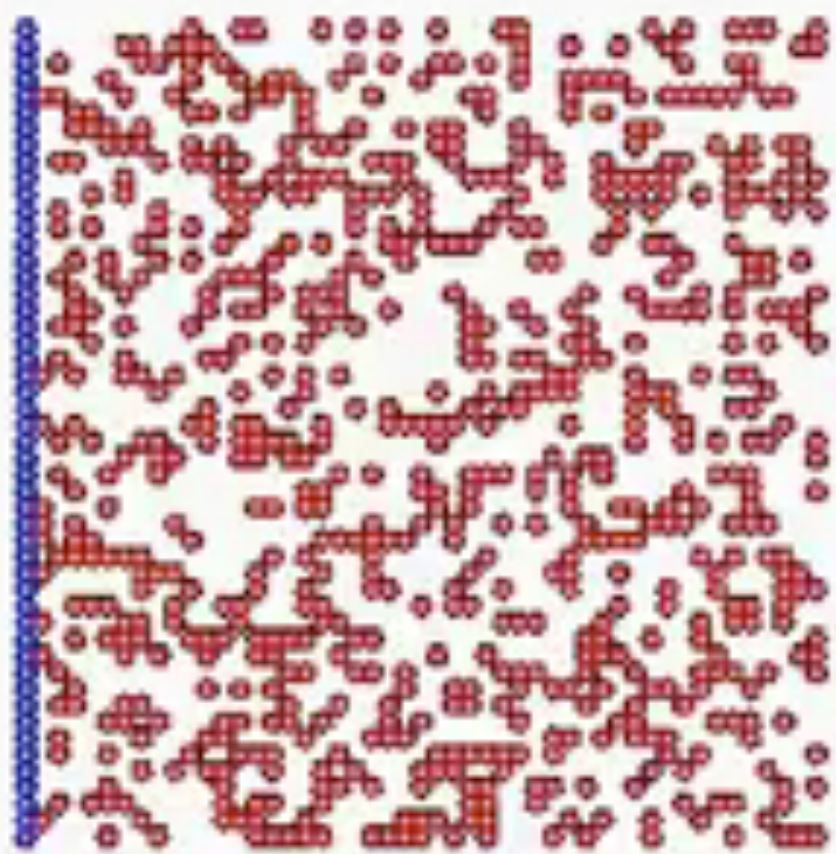
(a) initial lattice



(b) propagation



(c) collision handling





Liquid Simulation on Lattice-Based Tetrahedral Meshes

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