

Timing of Cyber Conflict

presented by Padraic Cashin

Robert Axelrod and Rumen Iliev

When do you attack?

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- When expending resources yields value greater than possible future value
- Each entity has a **Threshold**, T , for **Stakes**, s ; minimum level of stakes before an attack will be considered.
- Resources consist of exploits, back doors, bot nets, etc.

Model Assumptions

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Entities know the current stakes, but only know the distribution of future stakes
- Future stakes are out of your control
- Future effectiveness of a resource can only be estimated

Shelf Life of Resources

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Vulnerabilities can be discovered and patched.
- A vulnerability is **stealthy**, S, if it remains viable after use
- A vulnerability is **persistent**, P, if it remains viable when not used

Persistence vs Stealth

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Persistent resources are not currently deployed. Stealthy resources have already been used.
- $P = \Pr(\text{resource survives} \mid \text{not use it})$
- $S = \Pr(\text{resource survives} \mid \text{use it})$

Value vs Gain

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- The **gain**, G , of a resource is the immediate value from deploying a resource
- The **value**, V , of a resource is the sum of immediate gains and all future gains
- The value of a resource over time is discounted by a fixed percent, w

Defining Value

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

Value of a stealthy resource:

$$V_S = G(T) + wSV \quad (1)$$

Value of a persistent resource:

$$V_P = wPV \quad (2)$$

Expected value over-time:

$$V = \Pr(s \geq T)[G(T) + wSV] + (1 - \Pr(s \geq T))wPV \quad (3)$$

Determining Optimal Timing of Attacks

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Distribution of stakes is linear. Based on the role of a die.
- The discount rate is fixed at $w = 0.9$
- Analyse the effects of stealth and persistence on threshold

Effect of Persistence

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

T \ P	0.1	0.2	0.3	0.4	0.5	0.6	0.7	0.8	0.9	1.0
6	1.09	1.20	1.33	1.49	1.70	1.98	2.37	2.94	3.88	5.71
5	1.98	2.16	2.37	2.62	2.93	3.33	3.86	4.58	5.64	7.33
4	2.68	2.98	3.13	3.42	3.77	4.20	4.74	5.43	6.37	7.69
3	3.19	3.41	3.66	3.95	4.29	4.69	5.17	5.77	6.52	7.50
2	3.52	3.72	3.96	4.22	4.52	4.87	5.27	5.75	6.32	7.02
1	3.66	3.85	4.05	4.27	4.52	4.79	5.11	5.47	5.88	6.36

Stealth is set to half of Persistence. Using a resource doubles the likely hood it will be discovered.

Effect of Stealth

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

T \ S	0.1	0.2	0.3	0.4	0.5	0.6	0.7	0.8
6	2.60	2.70	2.82	2.94	3.08	3.23	3.39	3.57
5	3.74	3.99	4.26	4.58	4.95	5.39	5.91	6.55
4	4.20	4.55	4.95	5.43	6.02	6.76	7.69	8.93
3	4.29	4.69	5.17	5.77	6.52	7.50	8.82	10.71
2	4.14	4.57	5.09	5.75	6.60	7.75	9.39	11.90
1	3.85	4.27	4.79	5.47	6.36	7.61	9.46	12.50

Persistence is fixed at 0.8

Results

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- As Persistence increases Threshold increases
- As Stealth increases Threshold goes down
- Patience increases when stealth is low, persistence is high, and large stakes are rare

Case Study 1: Stuxnet

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Low Persistence, High Stealth, and High Stakes
- Multiple resources used at once, high cost of use
- Gain was not estimated properly due to source code leaks

Case Study 2: Attack on Saudi Aramco

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Broad attack on Saudi and US oil pipelines (30,000 workstations infected)
- Very High Stakes, Low Stealth
- Attackers immediately deployed a resource en masse

Case Study 3: Chinese Cyber Espionage

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Wide spread deployment of cyber resources
- Moderate Stealth against vigilant targets, Minimal Stakes
- Either persistence is very low or expect High Stealth against outliers

Case Study 4: Refusal to Export Minerals

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Chinese refused to export rare-earth minerals due to Japanese detainment of Chinese fishing crew.
- Very High Persistence, Low Stealth, Low Value
- China might have a artificially low threshold or low patience

Effect of Zero-Day Markets

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Increased pressure to find exploits leads to simultaneous discover; Decreases Persistence.
- Lower Persistence lowers Threshold; Increase resource deployment
- Prices predicted to drop as exploits become available and less persistent

Conclusions

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Model explains cyber conflict frequency using economic models
- Entities attempt to maximize resource effectiveness
- Resources are both perishable and detectable
- Assumes each entity will act perfectly

What Happens if an Attack is False?

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Evidence of an Attack can be Spoofed
- Attacks are not necessarily resource intensive; non-state attacks are possible
- Attacks can be attributed to different entities incorrectly, or can be left unattributed
- Each entity has unknown capabilities

The Blame Game

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- Two player Bayesian Game, players have imperfect knowledge of each other but can estimate a probability of state
- Players are either the Attacker (\mathcal{A}) or the Blamer (\mathcal{B})
- \mathcal{A} chooses to attack \mathcal{B} or not
- \mathcal{B} chooses to blame \mathcal{A} or not

Behavior and Equilibria

Timing of
Cyber Conflict

Robert
Axelrod and
Rumen Iliev

- \mathcal{A} attempts to determine if \mathcal{B} is **knowledgeable**
- \mathcal{B} attempts to determine if \mathcal{A} is **vulnerable**
- Equilibria exist if no attack – no blame or attack – blame occurs
- Third parties can disrupt cooperative equilibrium