

nano? REAL PROGRAMMERS USE emacs



HEY. REAL PROGRAMMERS USE vim.



WELL, REAL PROGRAMMERS USE ed.



NO, REAL PROGRAMMERS USE cat.



REAL PROGRAMMERS USE A MAGNETIZED NEEDLE AND A STEADY HAND.



EXCUSE ME, BUT REAL PROGRAMMERS USE BUTTERFLIES.



THEY OPEN THEIR HANDS AND LET THE DELICATE WINGS FLAP ONCE.

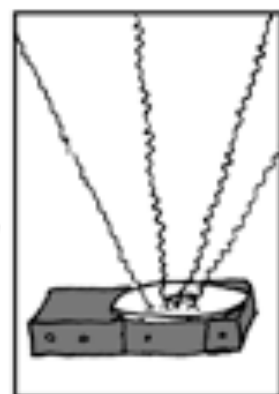
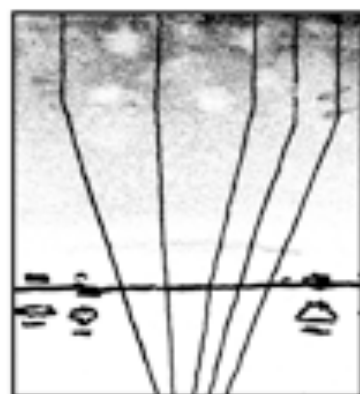


THE DISTURBANCE RIPPLES OUTWARD, CHANGING THE FLOW OF THE EDDY CURRENTS IN THE UPPER ATMOSPHERE.



THESE CAUSE MOMENTARY POCKETS OF HIGHER-PRESSURE AIR TO FORM,

WHICH ACT AS LENSES THAT DEFLECT INCOMING COSMIC RAYS, FOCUSING THEM TO STRIKE THE DRIVE PLATTER AND FLIP THE DESIRED BIT.



NICE.  
'COURSE, THERE'S AN EMACS COMMAND TO DO THAT.  
OH YEAH! GOOD OL' C-x M-c M-butterfly...



DAMMIT, EMACS.

# Logistics

- We plan to return graded projects on ~~Wednesday~~, **Thursday**.
- I will hold office hours all day Thursday so you can pick up and discuss your graded project.
- Project 2 has been posted.
- Groups have been assigned.

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¿Project Questions?

# Genetic Algorithms Continued

Lecture 10

# Terminology

**Amino acids** biochemical properties: nonpolar, polar, basic, acidic

Termination: stop codon

On chromosomes:

4 Bases {T,C,A,G}

Codons (sequences of 3 bases) code for each amino acid.

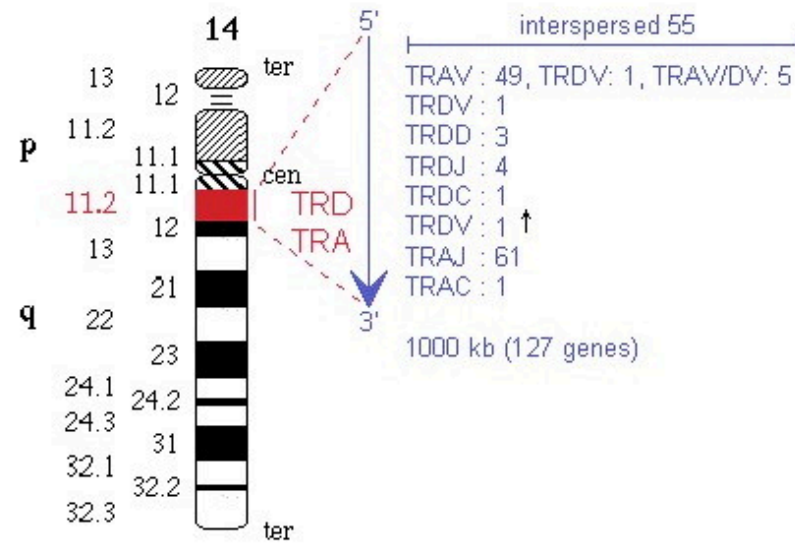
Genes (variable length sequences of codons) code for complete proteins.

Proteins form the Phenotype.

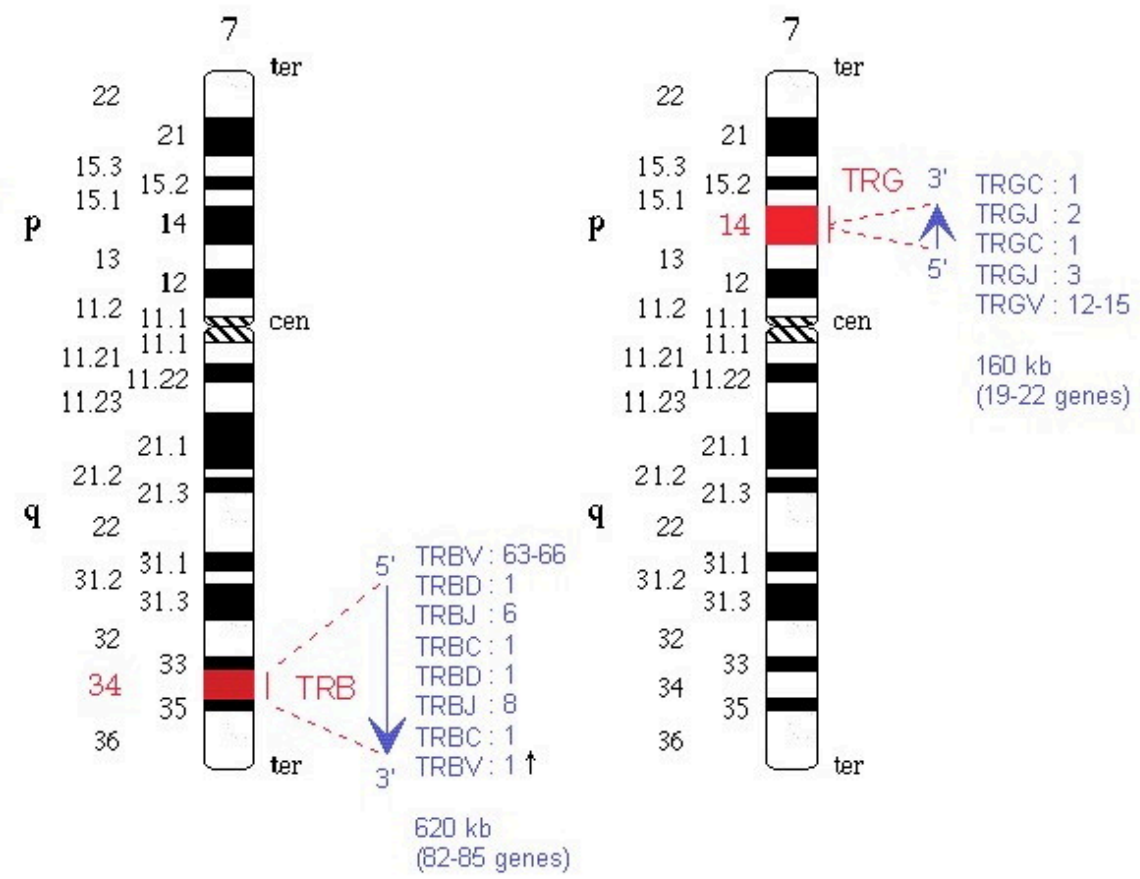
Standard genetic code

1st base	2nd base								3rd base
	T		C		A		G		
T	TTT	(Phe/F) Phenylalanine	TCT	(Ser/S) Serine	TAT	(Tyr/Y) Tyrosine	TGT	(Cys/C) Cysteine	T
	TTC		TCC		TAC		TGC		C
	TTA	(Leu/L) Leucine	TCA		TAA <sup>[B]</sup>	Stop (Ochre)	TGA <sup>[B]</sup>	Stop (Opal)	A
	TTG		TCG		TAG <sup>[B]</sup>	Stop (Amber)	TGG	(Trp/W) Tryptophan	G
C	CTT	(Leu/L) Leucine	CCT	(Pro/P) Proline	CAT	(His/H) Histidine	CGT	(Arg/R) Arginine	T
	CTC		CCC		CAC		CGC		C
	CTA		CCA		CAA	(Gln/Q) Glutamine	CGA		A
	CTG		CCG		CAG		CGG		G
A	ATT	(Ile/I) Isoleucine	ACT	(Thr/T) Threonine	AAT	(Asn/N) Asparagine	AGT	(Ser/S) Serine	T
	ATC		ACC		AAC		AGC		C
	ATA		ACA		AAA	(Lys/K) Lysine	AGA	(Arg/R) Arginine	A
	ATG <sup>[A]</sup>	(Met/M) Methionine	ACG		AAG		AGG		G
G	GTT	(Val/V) Valine	GCT	(Ala/A) Alanine	GAT	(Asp/D) Aspartic acid	GGT	(Gly/G) Glycine	T
	GTC		GCC		GAC		GGC		C
	GTA		GCA		GAA	(Glu/E) Glutamic acid	GGA		A
	GTG		GCG		GAG		GGG		G

# Chromosomal location of the TCR $\alpha/\delta$ , $\beta$ and $\gamma$ chain loci in man



locus representations from the International Immunogenetics Information System (IMGT) server (<http://www.imgt.org/>)



TCR: T cell Receptor

These loci are highly variable in their alleles.

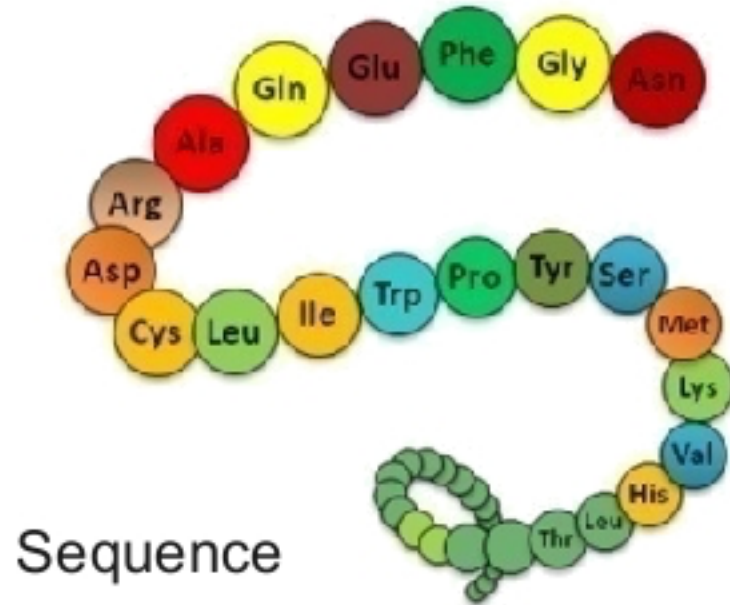
# Terminology: Protein Phenotype

Proteins form the phenotype.

The **conformation** of a protein is how it twists itself.

This conformation process is dynamical.

Proteins vibrate into low energy conformations.



© Robotics

Robot path planning techniques for protein folding. (Lydia Tapia, UNM)

Xinyu Tang, Shawna Thomas, Lydia Tapia, David P. Giedroc, Nancy M. Amato, "Simulating RNA Folding Kinetics on Approximated Energy Landscapes," *Journal of Molecular Biology*, 381(4):1055-1067, 2007

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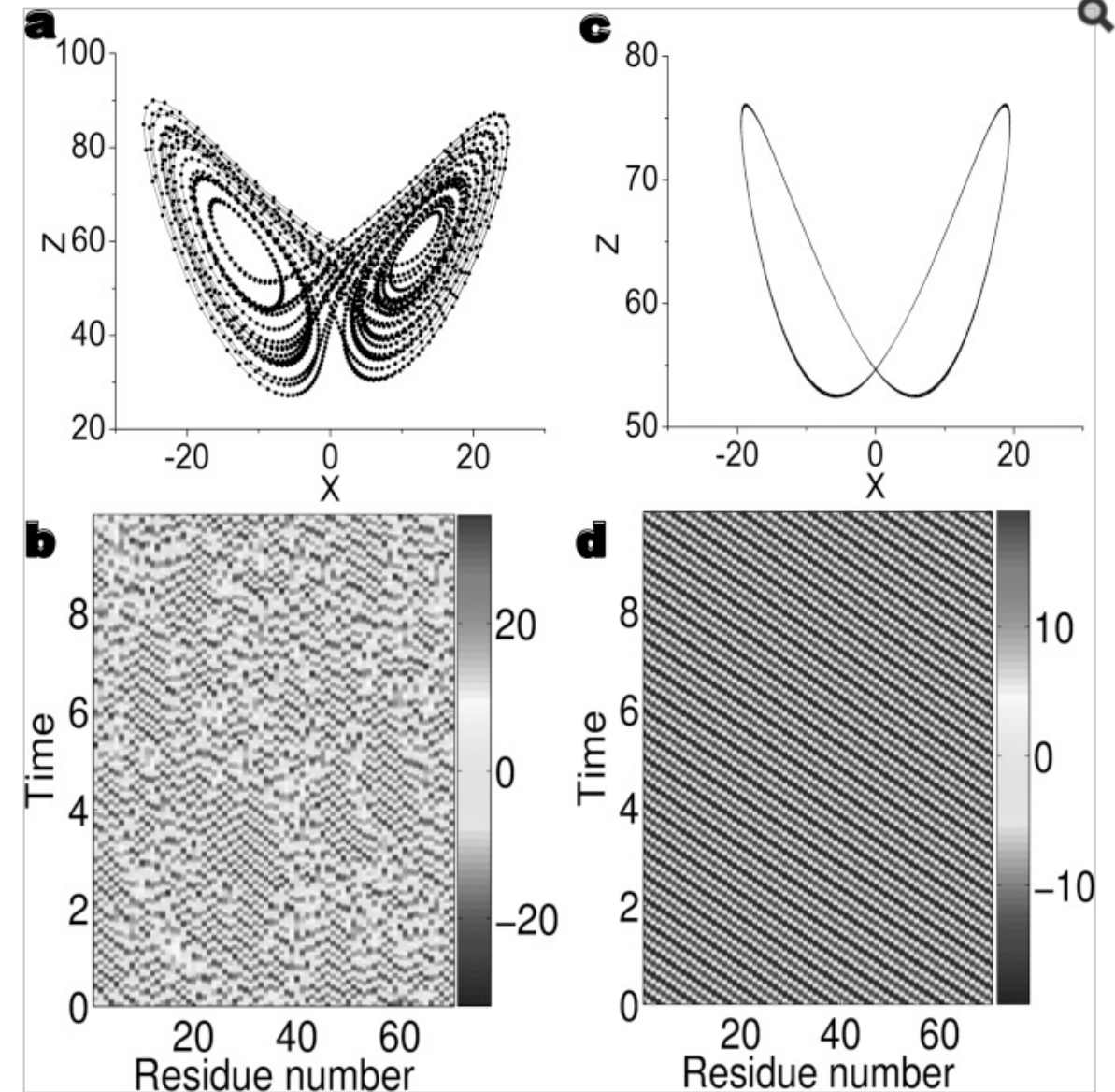
This conformation process is dynamical.

Proteins vibrate into low energy conformations.

These states are defined by stable and unstable fixed points and saddle points.

Xia, Kelin, and Guo-Wei Wei. "Molecular nonlinear dynamics and protein thermal uncertainty quantification." *Chaos: An Interdisciplinary Journal of Nonlinear Science* 24.1, 013103, 2014.

Figure 2

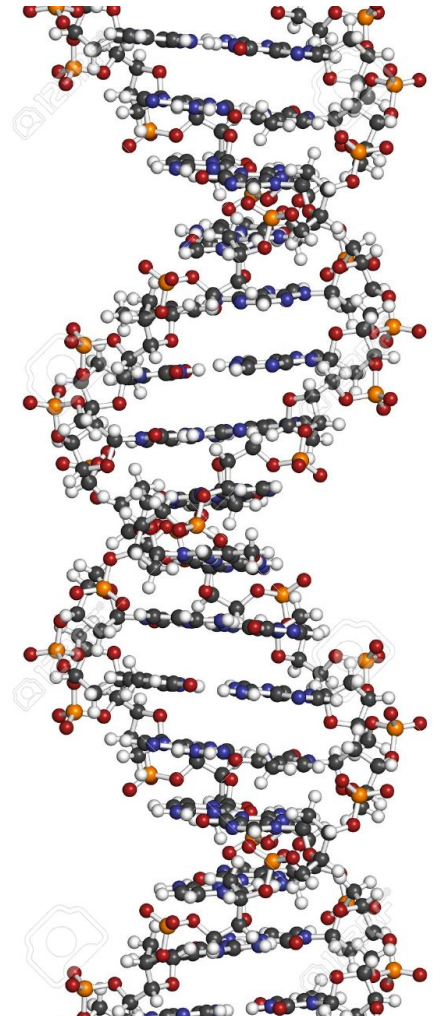


Transition from chaos to periodicity in the chaotic dynamics model (CDM) of bacteriocin AS-48 (PDB ID 1e68). (a) The butterfly wing pattern for one of 70 chaotic oscillators. (b) The solution of original 70 chaotic oscillators. (c) The periodic orbit of the ILDM for bacteriocin AS-48. (d) Bacteriocin AS-48 induced Hopf bifurcation from chaos. All of 70 nonlinear oscillators are in one lag synchronized periodic orbit.



# Terminology

- Gene – A variable in the genome
- Genotype – A string of symbols in the genome
- Phenotype – The decoding of a genome
- Locus – A position in the genome ( $i^{\text{th}}$  position in a string)
- Allele – A value the variable (gene) can take on.
- Epistasis – interdependence of genes (nonlinearity)
- Encoding – A mapping of a set of features into the genome.
- Biological genes can be independent of loci.



# Schemata

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Developing a theory of Genetic Algorithms



John Holland, Inventor of the GA.

The following slides are based on material from Goldberg, D. *Genetic Algorithms*, 1989 and Holland J, *Adaptation in Natural and Artificial Systems*, 1993

# GAs: Schemata

Consider a sample problem:

Encoding: integers as bitstrings.

Integers here are the **phenotype** and bitstrings are the **genotype**.

We want to maximise the bitstrings according to the fitness function:

$$F(x) = x^2$$

# GAs: Instructive Example

$$F(x) = x^2$$

<i>i</i>	Initial Pop.	Phenotype
1	01101	13
2	11000	24
3	01000	8
4	10011	19

# GAs: Instructive Example

$$F(x) = x^2$$

<i>i</i>	Initial Pop.	Phenotype	Fitness ( <i>f</i> )
1	01101	13	169
2	11000	24	576
3	01000	8	64
4	10011	19	361

# Roulette Selection

$$P(i) = \frac{f_i}{\sum_i f_i}$$

<i>i</i>	Initial Gen.	Phenotype	Fitness ( <i>f</i> )	<i>P</i> ( <i>i</i> )
1	01101	13	169	0.14
2	11000	24	576	0.49
3	01000	8	64	0.06
4	10011	19	361	0.31

After selection

$i$	2 <sup>nd</sup> Gen.	Phenotype	Fitness ( $f$ )
1	01101	13	169
2	11000	24	576
3	11000	24	576
4	10011	19	361

After selection, 1-point crossover

$i$	2 <sup>nd</sup> Gen.	Phenotype	Fitness ( $f$ )
1	0110 1	13	169
2	1100 0	24	576
3	11 000	24	576
4	10 011	19	361



After selection, crossover ( $P = 1.0$ ),

and mutation  $\left(P = \frac{1}{1000}\right)$

No mutation in this case.

$i$	2 <sup>nd</sup> Gen.	Phenotype	Fitness ( $f$ )
1	0110 0	12	144
2	1100 1	25	625
3	11 011	27	729
4	10 000	16	256

## GAs: Schemata

What information about the search space does this table contain?

Bitstring	Fitness
01101	169
11000	576
01000	64
10011	361

## GAs: Schemata

What information about the search space does this table contain?

Notice that bitstrings starting with 1s have higher fitness.

Bitstring	Fitness
01101	169
11000	576
01000	64
10011	361

## GAs: Schemata

1. Similarity among strings in the population.
2. A causal relationship between the strings and the fitness function.

Schemata capture this.

Bitstring	Fitness
01101	169
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Schemata define functional equivalence classes.

# GAs: Schemata

Schemata are defined over the string alphabet plus a metasymbol ‘\*’

‘\*’ is just a wildcard (not a Kleene star if you are in CS500).

# GAs: Schemata

Schemata (or similarity templates)  
define equivalence classes.

$10 * 01$

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$$10 * 01 = \{10101, 10001\}$$



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$$0 * 1 * *$$

# GAs: Schemata

Schemata (or similarity templates) define equivalence classes.

$$10 * 01 = \{10101, 10001\}$$

$$*000* = \{00000, 00001, 10000, 10001\}$$

$0 * 1 * *$  All strings of length 5 with a 0 in the first position and a 1 in the third.

# GAs: Schemata

How many schemata are there  
for an alphabet with cardinality  $k$   
and genome length  $N$ ?

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$$(k + 1)^N$$

# GAs: Schemata

We can bound the number of schemata in a population.

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In our example:

The number of schemata for an individual genome is  $2^5$ .

because it can take on its actual value or the wildcard.

(2 values)



# GAs: Schemata

We can bound the number of schemata in a population.

In our example:

The number of schemata for an individual genome is  $2^5$ .  
because it can take on its actual value or the wildcard.

If the population has  $n$  individuals  
there will be at most  $n2^2$  schemata.

# GAs: Schemata

So what? How does this help us?

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Can think in terms of useful classes of genomes, since some of the variation won't matter.

Holland makes an optimality argument for GAs using schemata.

# GAs: Schemata

So what? How does this help us?

Can think in terms of useful classes of genomes,  
since some of the variation won't matter.

# Schemata

Take a couple of minutes to discuss with your neighbour the effect on Schemata in the current population of:

- **Reproduction**
- **Crossover**
- **Mutation**



# GAs: Schemata

First we consider reproduction:

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Schemata with higher fitness tends to increase each generation. No new schemata appear.

# GAs: Schemata: Reproduction

First we consider reproduction:

Schemata with higher fitness tends to increase each generation. No new schemata appear.

Formally, for a schema  $S$  in population  $A$ ,

The number of individuals with schema  $S$  at time  $t + 1$ , given the number at time  $t$  is

$$m_S(t + 1) = m_S(t) \frac{f(S)}{\text{mean}(f(A))}$$



# GAs: Schemata: Reproduction

Consider a schema with fitness  $c \times \text{mean}(A)$ , and  $c > 1$

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We can rewrite the increase in representation for this schema,  $S$ , as:

$$m_S(t + 1) = m_S(t) \left[ \frac{c \times \text{mean}(A)}{\text{mean}(A)} \right]$$

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$$m_S(t) = c \times c \times \cdots \times c \times m_S(0)$$

# GAs: Schemata: Reproduction

We can rewrite the increase in representation for this schema,  $S$ , as:

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$$m_S(t + 1) = c \times m_S(t)$$

$$m_S(t) = c \times c \times \cdots \times c \times m_S(0) = c^t \quad \text{Exponential}$$

# GAs: Schemata

Now consider the effect of 1-point crossover on schemata.

## GAs: Schemata

Now consider the effect of 1-point crossover on schemata.

Schemata survive if they are not cut by the crossover

For example:


1 \* \* \* 0 is ~~less~~<sup>more</sup> likely than \* \* \* 10\*

to be destroyed.

Formally:

Probability of disruption:  $P_D = P_{\times} \frac{D(S)}{L_S - 1}$  where,

$P_{\times}$  is the crossover probability,

 1 fewer crossover sites than genes

$D(S)$  is the defining length of schema  $S$ ,

$L$  is the string/genome length.

Defining length: distance between the first and last non-wildcard symbols.



GAs: Schemata

How about mutation?

## GAs: Schemata

Mutation is more likely to destroy "higher order" schemata.  
The order is the number of fixed (non-wildcard) symbols.

# GAs: Schemata

Mutation is more likely to destroy "higher order" schemata.  
The order is the number of fixed (non-wildcard) symbols.

**More formally,**

Probability of disruption:  $P_D = \left(1 - (1 - P_{\text{mut}})^{O(S)}\right)$  where,

$P_{\text{mut}}$  is the mutation probability,

$O(S)$  is the order of the schema  $S$ .

GAs: Schemata

**Also called the fundamental  
Theorem of Genetic Algorithms**

The schema theorem:

Fit schema with lower defining length  
and lower order increase exponentially in the  
population over time.

These are called “building blocks”

# GAs: Schemata

**Also called the fundamental  
Theorem of Genetic Algorithms**

## Schema Theorem

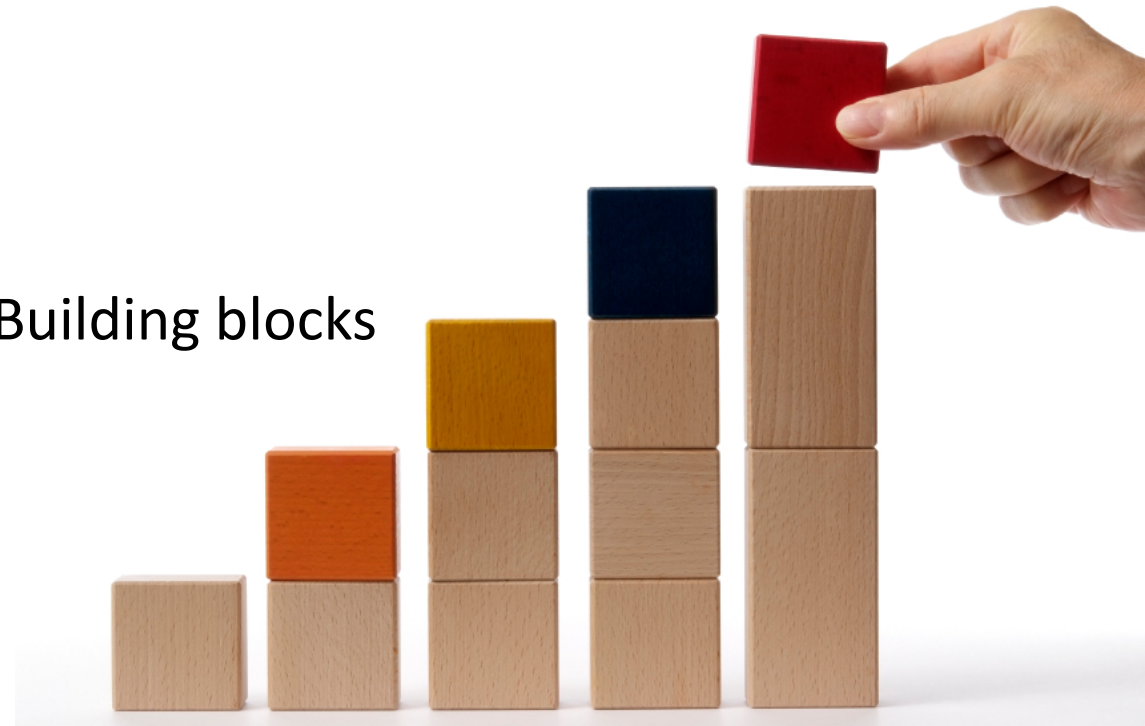
More formally,

$$m_S(t + 1) \geq m_S(t) \frac{f(S)}{\text{mean}(S)} \left[ 1 - P_{\times} \frac{D(S)}{L - 1} - O(S) P_{\text{mut}} \right]$$

# GAs: Schemata

Building block hypothesis: A genetic algorithm seeks optimal performance through the juxtaposition of short, low-order, high- performance schemata, called the building blocks.

Premature convergence is the main challenge. Building blocks can cause premature convergence.



# A Royal Road Problem: Testing the Building Block Hypothesis

$s_1 = 11111111*****$ ;  $c_1 = 8$   
 $s_2 = *****11111111*****$ ;  $c_2 = 8$   
 $s_3 = *****11111111*****$ ;  $c_3 = 8$   
 $s_4 = *****11111111*****$ ;  $c_4 = 8$   
 $s_5 = *****11111111*****$ ;  $c_5 = 8$   
 $s_6 = *****11111111*****$ ;  $c_6 = 8$   
 $s_7 = *****11111111*****$ ;  $c_7 = 8$   
 $s_8 = *****11111111$ ;  $c_8 = 8$   
 $s_9 = 1111111111111111*****$ ;  $c_9 = 16$   
 $s_{10} = *****1111111111111111*****$ ;  $c_{10} = 16$   
 $s_{11} = *****1111111111111111*****$ ;  $c_{11} = 16$   
 $s_{12} = *****1111111111111111$ ;  $c_{12} = 16$   
 $s_{13} = 11111111111111111111111111111111*****$ ;  $c_{13} = 32$   
 $s_{14} = *****11111111111111111111111111111111$ ;  $c_{14} = 32$   
 $s_{15} = 11$ ;  $c_{15} = 64$

Figure 1: Example Royal Road Function.  $F(x) = \sum_{s \in S} c_s \sigma_s(x)$ , where  $x$  is a bit string,  $c_s$  is a value assigned to the schema  $s$  (here,  $c_s = order(s)$ ), and  $\sigma_s(x) = \begin{cases} 1 & \text{if } x \text{ is an instance of } s \\ 0 & \text{otherwise.} \end{cases}$

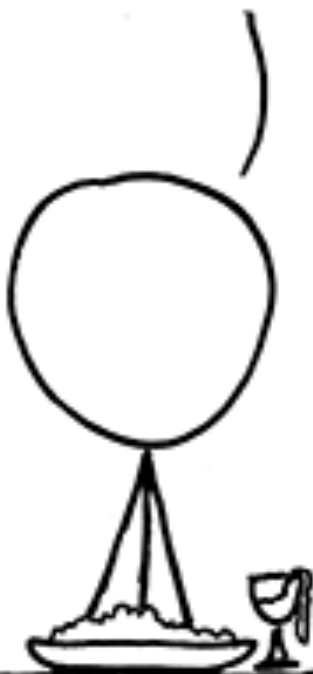
Mitchell, Melanie, Stephanie Forrest, and John H. Holland. "The royal road for genetic algorithms: Fitness landscapes and GA performance." *Proceedings of the first european conference on artificial life*. 1992.

I'VE GOT...  
CHEERIOS  
WITH A SHOT  
OF VERMOUTH.

AT LEAST IT'S BETTER  
THAN THE QUAIL EGGS  
IN WHIPPED CREAM AND  
MSG FROM LAST TIME.

ARE THESE  
SKITTLES  
*DEEP-FRIED?*

C'MON, GUYS, BE PATIENT. IN A  
FEW HUNDRED MORE MEALS, THE  
GENETIC ALGORITHM SHOULD CATCH  
UP TO EXISTING RECIPES AND START  
TO OPTIMIZE.



WE'VE DECIDED TO DROP THE CS DEPARTMENT FROM OUR WEEKLY DINNER PARTY HOSTING ROTATION.



# Logistics

- Midterm exam - March 10<sup>th</sup>.
- Midterm review moved to March 6<sup>th</sup>.
  
- Project 2 is due on March 10<sup>th</sup>.
- Project reviews are due on March 20<sup>th</sup>.
- The in class competition will be on March 20<sup>th</sup>.
  
- Transitioning from Genetic Algorithms to Cellular Automata
- We will be using genetic algorithms in the next project to explore cellular automata and game theory.

# Theory of Self-Reproducing Automata

## - John von Neumann

- David Shubsda
- Joshua Ridents
  
- Turn in your review forms in to Bianca at the end of the presentation.
- I will put David and Joshua's slides on the course website.

# GAs: k-armed bandits

Suppose you can play  $k$  slot machines.

The  $i^{\text{th}}$  machine pays jackpot following a

Gaussian probability distribution with

mean  $= \mu_i$ , and variance  $= \sigma_i^2$



# GAs: k-armed bandits

Suppose you can play  $k$  slot machines.

The  $i^{\text{th}}$  machine pays jackpot following a

Gaussian probability distribution with

mean  $= \mu_i$ , and variance  $= \sigma_i^2$



The challenge is to win the most money possible (or lose the least) over time.

GAs: k-armed bandits

Reason about the best trade-off



# GAs: k-armed bandits

Reason about the best trade-off

A strategy:

Given a maximum of  $N$  pulls at the arm we could try:

Allocating an equal number,  $n$ , of pulls, where  $(kn < N)$ , to each arm, then use all the remaining time to pull on the arm with the most payouts.



# GAs: k-armed bandits

Reason about the best trade-off



Given  $N$ ,  $\mu_i$ , and  $\sigma_i^2$ ,  
we can define a loss function,  $L(N, n)$ .

# GAs: Optimality of Trial Allocation

- There is a trade-off between sampling near the best observed and exploring the fitness landscape.
- Loss due to searching near the currently known optimum is due to *sampling error*.
- Loss due to choosing areas that are not known to be good we might call *performance loss*.



# GAs: Optimality of Trial Allocation

- The presence of a trade-off suggests an optimisation problem.

# GAs: k-armed bandits

Reason about the best trade-off

Given  $N$ ,  $\mu_i$ , and  $\sigma_i^2$ ,

we can define a loss function,  $L(N, n)$ .

$$L(N, n) = |\mu_1 - \mu_2| [(N - n)q(n) + n(1 - q(n))]$$

The 2-arm case generalises



# GAs: k-armed bandits

$$L(N, n) = |\mu_1 - \mu_2| [(N - n)q(n) + n(1 - q(n))]$$

Where  $q(n)$  is the probability that the worst arm is the best arm after  $n$  pulls.

In other words  $q(n)$  is the probability that you got unlucky as a function of the explore/exploit trade-off.

$q(n)$  decreases exponentially with  $n$ .  
 To minimise  $q(n)$  we allocate  $O(e^n)$  trials to the known best.

“The ratio of trials of the observed best to the second best grows exponentially.”

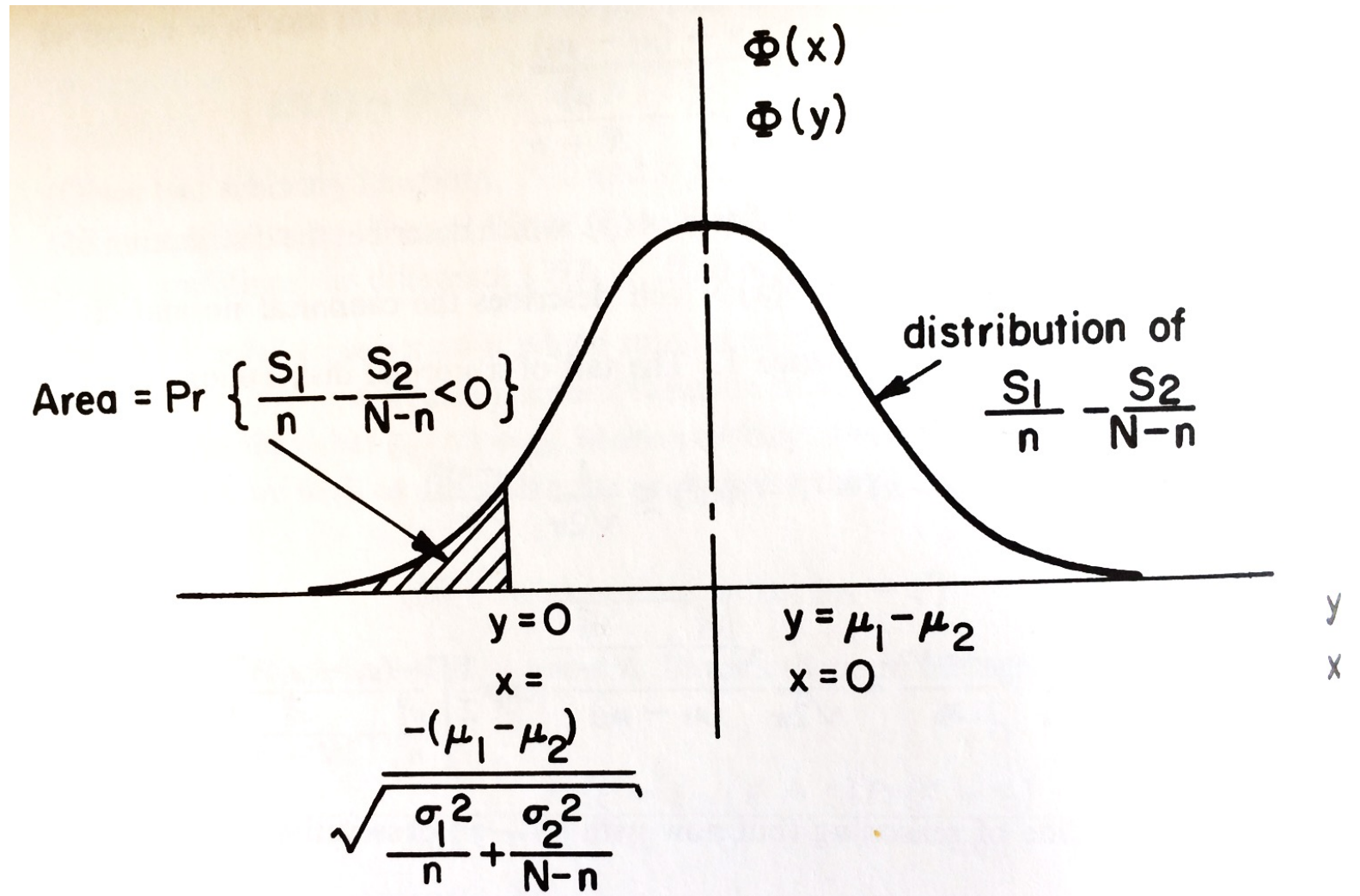


Fig. 11. The convolution of  $\frac{S_1}{n}$  with  $\frac{S_2}{N-n}$

# GAs: k-armed bandits

Holland showed that allocating an exponential number of pulls to the best performing slot machine minimises the loss function  $L(N, n)$ .

This exponential allocation to exploitation best solves the problem.

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**What does this have to  
Do with the Fundamental Theorem?**

# GAs: k-armed bandits

Holland showed that allocating an exponential number of pulls to the best performing slot machine minimises the loss function  $L(N, n)$ .

This exponential allocation to exploitation best solves the problem.

**What does this have to  
Do with the Fundamental Theorem?**

**The fundamental theorem shows that GAs  
allocate exponentially increasing samples to known best solutions.**

# Diversity

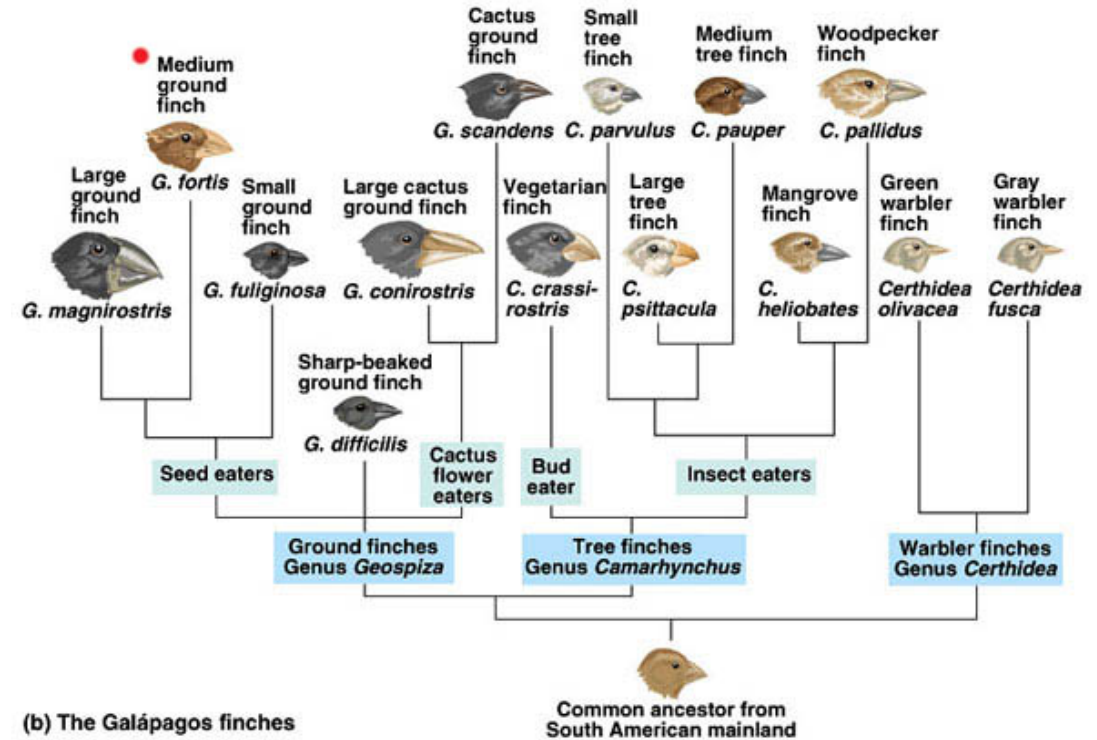
- A pitfall of GAs is that they may converge too early on a local maxima.
- If the population were infinite this wouldn't be a problem.
- The smaller the population the more likely they are to converge prematurely.
- Building blocks can increase premature convergence for some problems.



# Schemata: Implicit Parallelism

Holland also showed in 1985 that if  $K$  is the number of strings processed each generation and  $N$  is the number of schemata  
Then  $N \in O(K^3)$ .

# Diversity: Island GAs



(b) The Galápagos finches

Copyright © Pearson Education, Inc., publishing as Benjamin Cummings.

Recall Darwin's famous finches. The diversity is in part due to evolution on islands.

# Diversity: Island GAs

What would this look like as an algorithm?

# Diversity: Elitism

- Diversity is great for avoiding local minima.
- How do we keep from moving too far from regions we know are good.
- Roulette selection can easily discard the most fit individuals.

# Diversity: Elitism

- Diversity is great for avoiding local minima.
- How do we keep the population from forgetting about regions we know are good.
- Roulette selection can easily discard the most fit individuals.
- High mutation rates can move the whole population to a lower local maxima.
- Copying the highest fitness individual into the next generation unchanged is called *elitism*.

# Diversity: Elitism

- Elitism guarantees that the GA will converge\*

(Applies to simulated annealing and artificial immune system optimization as well)

\*Villalobos-Arias et al, “Asymptotic Convergence of Some Metaheuristics Used for Multiobjective Optimization”, *Foundations of Genetic Algorithms*, 2005

# Selection Pressure: Tournaments

- Selection pressure is a measure of how harsh we make the world.
- At one extreme only the very fittest individual would survive.
- At the other extreme everyone survives.

# Selection Pressure: Tournaments

- If selection pressure is too high the GA will converge prematurely on a local minima.
- If selection pressure is too low the GA will not converge at all.



# Selection Pressure: Tournaments

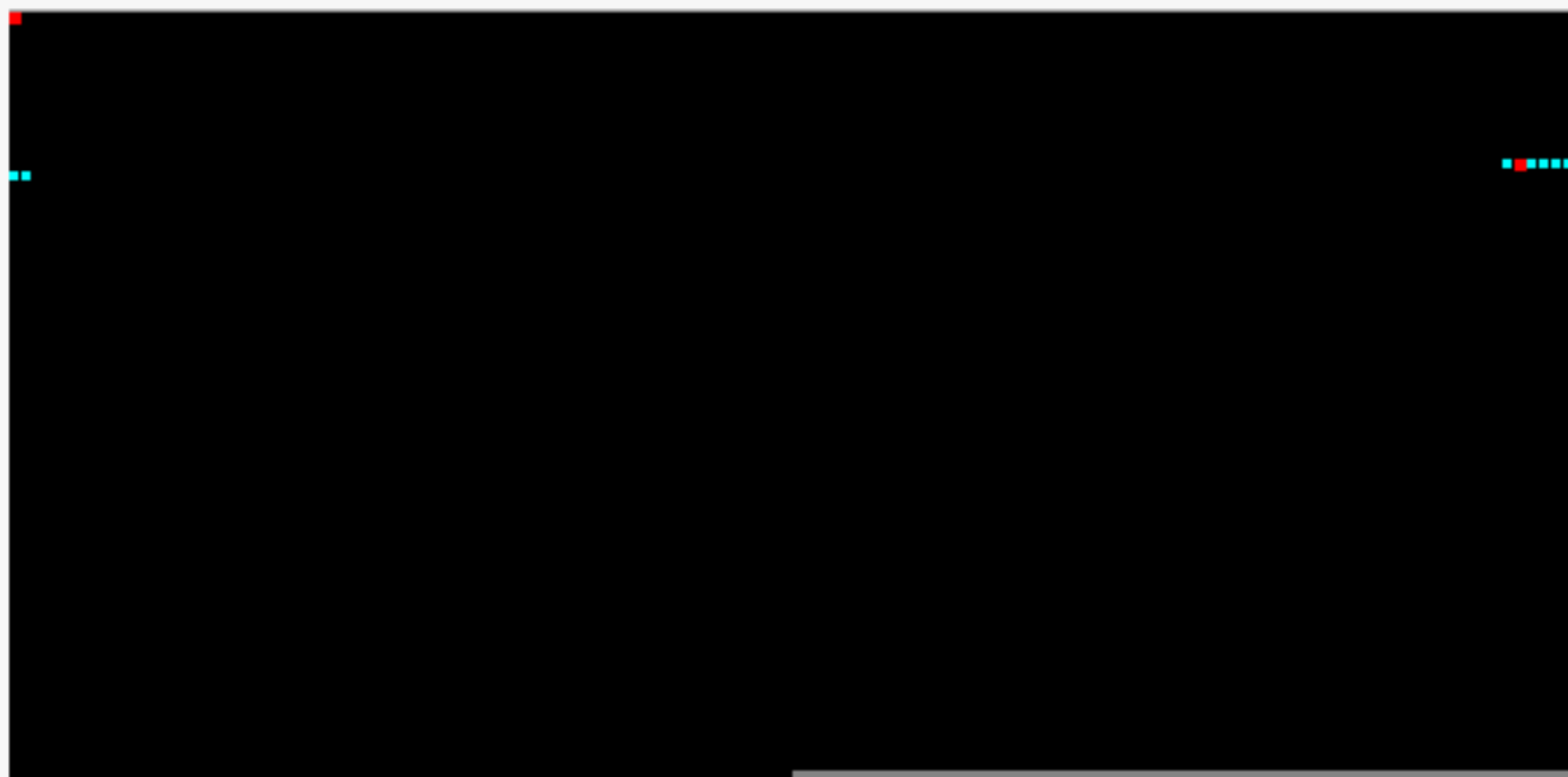
- Tournament selection is one way to tune the selection pressure.

# Evolving Programs

Project 2

# Project 2: Core Wars

- Alexander Dewdney, Mathematician and Computer Science
- Wrote a series of articles called Mathematical Recreation for Scientific American
- One of these was Core Wars
- Inspired by a real life incident (Creeper and Reaper).
- Christopher Langton invited Dewdney to present core wars  
At the first Alife conference.
- Playing with Core Wars was once considered dangerous.
- Capable of self-mutation...



Battle Info

Core Size: 8000      Max Cycles: 80000

Battle Status: Stopped

Completion:  0

Tournament Type: Round Robin

Round 1 of 1      Matchup 1 of 6

Max Processes: 8000

<input type="button" value="Imp"/>	<input type="text" value="0%"/>	1
<input type="button" value="Mice"/>	<input type="text" value="0%"/>	1
<input type="button" value="Midget"/>	<input type="text" value=""/>	
<input type="button" value="Piper"/>	<input type="text" value=""/>	

# Some x86 Assembly language

Address	Instruction	Register or RAM Address
77E814EE	mov	esi,dword ptr [edi+8]
77E814F1	mov	dword ptr [ebp+64h], 0Ah
77E814F8	add	esi, 4Ah
77E814FB	jmp	77E7E91A
77E81500	push	ebx
77E81501	xor	ebx,ebx
77E81503	cmp	ecx,ebx
77E81505	push	esi
77E81506	push	edi

```
;redcode
```

```
;author: T77
```

```
;name: Example1
```

```
;assert CORESIZE=8000 && MAXLENGTH > 100
```

```
MOV 0, 1
```

# - immediate addressing

;redcode

;author: T77

;name: Example2

;assert CORESIZE=8000 && MAXLENGTH > 100

ADD #4, 3 ← Program counter is here

MOV 2, @2

JMP -2

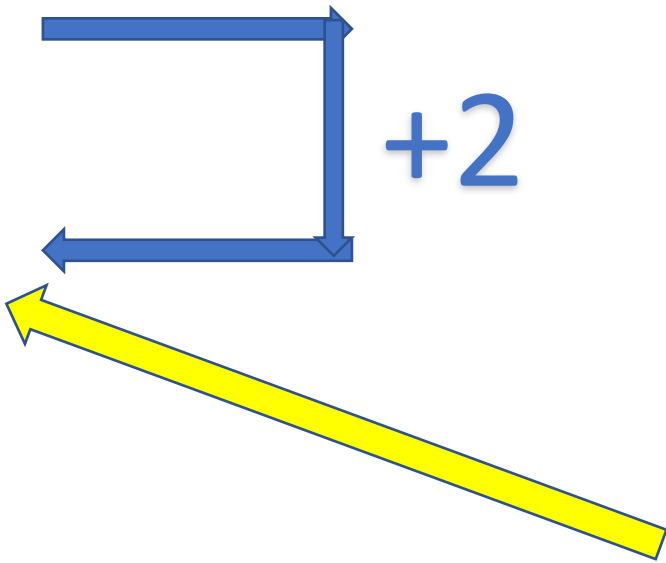
DAT #0, #0

```
;redcode  
;author: T77  
;name: Example2  
;assert CORESIZE=8000 && MAXLENGTH > 100  
ADD #4, 3  
MOV 2, @2 ← Program counter is here  
JMP -2  
DAT #0, #4
```



# @ - indirect addressing

```
;redcode  
;author: T77  
;name: Example2  
;assert CORESIZE=8000 && MAXLENGTH > 100  
ADD #4, 3  
MOV 2, @2  
JMP -2  
DAT #0, #4  
.br/>.br/>.br/>DAT #0, #4
```



B field points here

# @ is indirect addressing

```
;redcode
```

```
;author: T77
```

```
;name: Example2
```

```
;assert CORESIZE=8000 && MAXLENGTH > 100
```

```
ADD #4, 3
```

```
MOV 2, @2
```

```
JMP -2
```

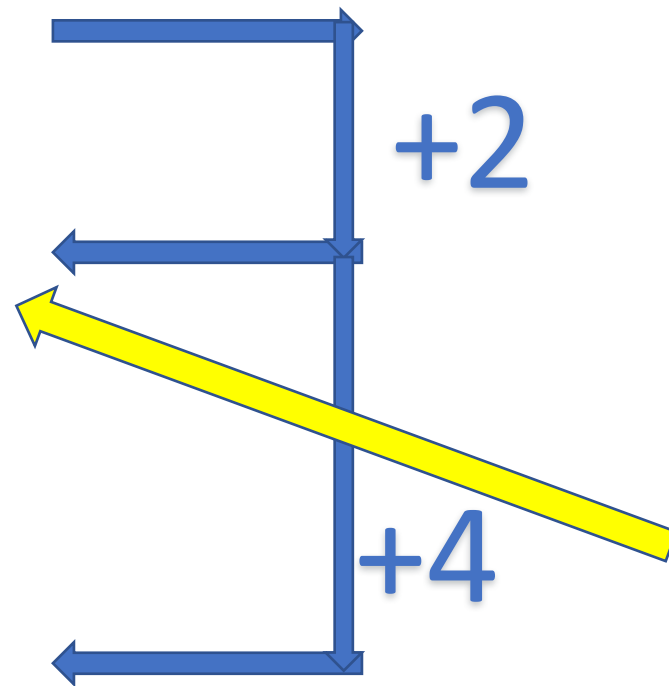
```
DAT #0, #4
```

```
.
```

```
.
```

```
.
```

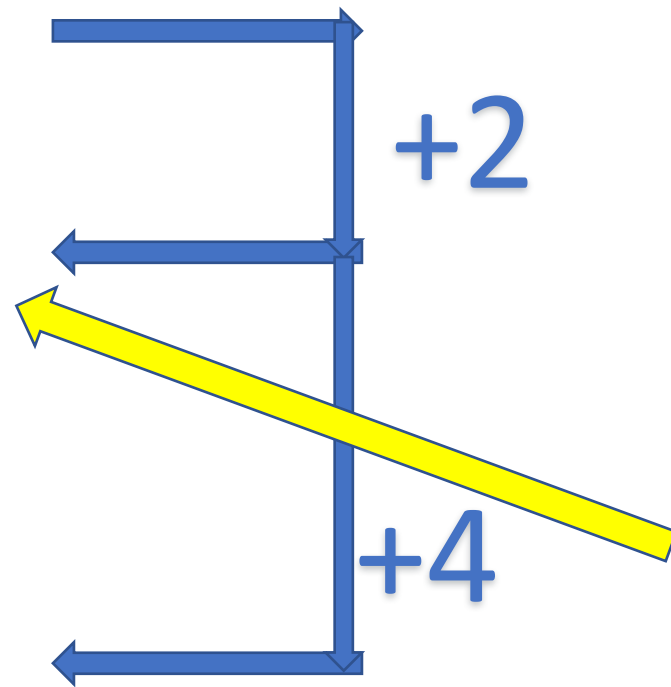
```
DAT #0, #4
```



B field points here

This program places data separated by 4 addresses forever.

```
ADD #4, 3  
MOV 2, @2  
JMP -2  
DAT #0, #4  
.br/>.br/>.br/>DAT #0, #4
```



B field points here

```
;redcode
```

```
;author: T77
```

```
;name: Example3
```

```
;assert CORESIZE=8000 && MAXLENGTH > 100
```

```
spl 0
```



Create new thread

```
mov 2, <-1
```

And continue with next

```
jmp -1, -1
```

instruction

```
;redcode
```

```
;author: T77
```

```
;name: Example3
```

```
;assert CORESIZE=8000 && MAXLENGTH > 100
```

```
spl 0
```

```
mov 2, <-1
```

```
jmp -1, -1
```

< Indirect with  
predecrement

